EE582

Physical Design Automation of VLSI Circuits and Systems

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What We Will Study

- Interconnect analysis
 - Elmore delay
 - Simple timing analysis
- Interconnect Optimization
 - Dynamic-programming-based buffer insertion
 - Two-pin nets
 - Multi-pin nets
 - Single buffer type
 - Multiple buffer types



Delay calculation (Elmore delay)





Delay calculation (Elmore delay)

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Timing analysis



- CLK period: T_{CLK}
- Path delay: $d_p = nd_1 + gd_1 + nd_2 + gd_2 + nd_3 + gd_3 + nd_4$
- Setup time: t_s

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- Required Arrival Time (RAT) = $T_{CLK} t_s$
- Arrival Time (AT) = Signal delay = Path delay = d_p

- Slack = RAT - AT =
$$T_{CLK} - t_s - d_p$$

Timing analysis



- Required Arrival Time (RAT) = $T_{CLK} - t_s$

RAT = 900ps AT = 800ps Slk = +100ps

- Arrival Time (AT) = Signal delay = Path delay = d_p
- Slack = RAT AT = $T_{CLK} t_s d_p$
- Example

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- T_{CLK}: 1ns
- d_p: 800ps
- t_s: 100ps
- Slack: 1ns (800ps) 100ps = +100ps

- Timing analysis
 - Positive slack: We have some margin for optimization.
 - Negative slack: We are violating the timing constraint.





- How to optimize interconnects
 - Gate sizing
 - Wire sizing
 - Buffer insertion
 - Placement adjustment
 - Gate decomposition
 - New routing topology generation



- Problem definition for buffer insertion
 - Given
 - An RC tree or a routing topology
 - Characteristics (R_{OUT}) of the driver node
 - Characteristics (C_{LOAD}) of the sink nodes
 - Characteristics (R_{OUT} and C_{IN}) of a buffer type
 - Characteristics (r_{wire} and c_{wire}) of the wires
 - Timing constraints (RAT at each sink node)
 - Bufferable locations
 - Maximize
 - Slack at the driver node











• Example

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- $r_{wire} = 2\Omega/um$
- $c_{wire} = 0.4 fF/um$
- $R_{OUT} = 2k\Omega$
- $C_{IN} = 5 fF$
- D_{BUF} = 50ps



Delay = 2k * (805f) + 0.4k * (765f) + 0.4k * (685f) + 0.4k * (605f) + ... + 0.4k * (125f) + 0.4k * (45f) = 3230ps



• Example

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Delay = 2k * (405f) + 0.4k * (365f) + 0.4k * (285f) + 0.4k * (205f) + 0.4k * (125f) + 0.4k * (45f) + 50ps (buffer delay)

- + 2k * (405f) + 0.4k * (365f) + 0.4k * (285f) + 0.4k * (205f) + 0.4k * (125f) + 0.4k * (45f)
- = 1220ps + 50ps + 1220ps = 2490ps







- Dynamic programming-based buffer insertion
 - van Ginneken algorithm





- Dynamic programming-based buffer insertion
 - van Ginneken algorithm





Comparison of two candidates

RAT	Сар	Action
RAT (C_1) > RAT (C_2)	Cap (C_1) > Cap (C_2)	Keep both
RAT (C_1) > RAT (C_2)	Cap (C_1) < Cap (C_2)	Keep C ₁ only
RAT (C_1) < RAT (C_2)	Cap (C_1) > Cap (C_2)	Keep C ₂ only
RAT (C_1) < RAT (C_2)	Cap (C_1) < Cap (C_2)	Keep both



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How many candidates do we have at each location?





After candidate construction





• Example and back-traversal for the construction of the final solution





• Example





Multi-fanout nets





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• Multi-fanout nets

- (RAT, Cap) of $(C_{1,1} + C_{2,1}) = (MIN(RAT_{1,1}, RAT_{2,1}), Cap_{1,1} + Cap_{2,1})$
- (RAT, Cap) of $(C_{1,1} + C_{2,2}) = (MIN(RAT_{1,1}, RAT_{2,2}), Cap_{1,1} + Cap_{2,2})$

- ...



Branch 2 Cand 1 Cand 2 . . . Cand k Branch 1



Multi-fanout nets

Multi-fanout nets





- # candidates at a branch point
 - a candidates from branch 1
 - b candidates from branch 2







- # candidates at a branch point
 - a candidates from branch 1
 - b candidates from branch 2

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 C_1 (100, $Cap_{1,1} + Cap_{2,1}$) C_2 (100, $Cap_{11} + Cap_{22}$) Branch 1 Branch 2 C_4 (100, $Cap_{11} + Cap_{24}$) C_3 (100, $Cap_{11} + Cap_{23}$) C₄ (100, Cap_{1 1} + Cap_{2.4}) C_{11} (100, Cap_{11}) C_{21} (180, Cap_{21}) $C_{2,2}$ (160, $Cap_{2,2}$) C_{12} (80, Cap_{12}) Х = | $C_{2,3}$ (140, $Cap_{2,3}^{-,-}$) $C_{1,3}$ (60, $Cap_{1,3}$) C_5 (80, Cap₁₂ + Cap₂₁) C₁₄ (40, Cap₁₄) C₂₄ (120, Cap₂₄) C_{6} (80, $Cap_{1,2} + Cap_{2,2}$) C_{8} (80, $Cap_{12} + Cap_{24}$) $C_7 (80, Cap_{12} + Cap_{23})$ C_{8} (80, $Cap_{12} + Cap_{24}$) Formula C_9 (60, $Cap_{1,3} + Cap_{2,1}$) **RAT'** = MIN (RAT₁, RAT₂) C_{10} (60, $Cap_{13} + Cap_{22}$) C_{12} (60, $Cap_{13} + Cap_{24}$) $Cap' = Cap_1 + Cap_2$ C_{11} (60, $Cap_{13} + Cap_{23}$) C_{12} (60, $Cap_{13} + Cap_{24}$) C_{13} (40, $Cap_{14} + Cap_{21}$) $Cap_{11} > Cap_{12} > ... > Cap_{1a}$ C_{14} (40, $Cap_{14} + Cap_{22}$) C_{16} (40, $Cap_{1,4} + Cap_{2,4}$) C_{15} (40, $Cap_{14} + Cap_{23}$) C_{16} (40, $Cap_{14} + Cap_{24}$) WASHINGTON STATE

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Four candidates survive.

- # candidates at a branch point
 - a candidates from branch 1
 - b candidates from branch 2

Seven candidates survive.





candidates at a branch point (max. a + b – 1)
– RAT





candidates: O(n) where n is # bufferable locations





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- Problem definition for buffer insertion with different buffer types
 - Given
 - An RC tree or a routing topology
 - Characteristics (R_{OUT}) of the driver node
 - Characteristics (C_{LOAD}) of the sink nodes
 - Characteristics (R_{OUT} and C_{IN}) of each buffer type
 - Characteristics (r_{wire} and c_{wire}) of the wires
 - Timing constraints (RAT at each sink node)
 - Bufferable locations
 - Maximize

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• Slack at the driver node

• Example





- Handling multiple buffer types can be done in the same way.
- The complexity goes up.



- Dynamic programming-based buffer insertion
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- # candidates (B is # buffer types)
 - 1st position
 - B+1 (un-buffered, $b_1, ..., b_B$)
 - 2nd position
 - (B+1)*(B+1) B*B = 2B+1
 - 3rd position
 - $(2B+1)^*(B+1) (2B)^*B = 3B+1$
 - 4th position
 - $(3B+1)^*(B+1) (3B)^*B = 4B+1$

— ...

candidates: O(B)

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More issues – Over-optimization
– DP is for each net.





More issues – Over-optimization
– DP is for each net.





- More issues Over-optimization
 - DP maximizes the RAT at the source node.
 - i.e., DP minimizes the delay from the source to the worst sink.
 - Problems when the timing conditions of the sinks are unbalanced.









