Homework Assignment 9

(Due May 3rd, 12pm)

- 1. [Carry Select Adder, **40 points**] In the lecture note, we used k-bit adders to design an Nbit carry select adder. However, we can use variable-length adders instead of fixed-length adders. The following shows the specification of logic blocks we are going to use:
 - N: 64
 - Delay of a full adder: $\Delta_{FA} = 100 ps$
 - Delay of a k-bit ripple carry adder (RCA): $k \cdot \Delta_{FA}$
 - Delay of a k-bit MUX (when $k \ge 10$) and a CO logic: $\varepsilon = 150ps$
 - Architecture: We split N into four groups as follows:



Since N is 64, the above architecture should satisfy the following equation:

$$k_1 + k_2 + k_3 + k_4 = 64$$

We also assume that $k_i \ge 10$ ($i = 1 \sim 4$).

- 1) Represent delay of the 64-bit adder shown above as a function of k_1, k_2, k_3, k_4 , and the MAX(a,b) function where "MAX(a, b) = a if (a>b) or b (if b>a)".
 - Delay of the k₁-bit RCA: $d_1 = k_1 \cdot \Delta_{FA} = 100k_1$
 - Delay of the k₂-bit RCA: $d_2 = k_2 \cdot \Delta_{FA} = 100k_2$
 - Delay of the k₃-bit RCA: $d_3 = k_3 \cdot \Delta_{FA} = 100k_3$
 - Delay of the k₄-bit RCA: $d_4 = k_4 \cdot \Delta_{FA} = 100k_4$
 - Arrival time (AT) at the output of the k₂-bit MUX and the first Carry-Out logic: $d_5 = MAX(d_1, d_2) + \varepsilon = 150 + MAX(100k_1, 100k_2)$
 - AT at the output of the k₃-bit MUX and the second Carry-Out logic: $d_6 = MAX(d_3, d_5) + \varepsilon = 150 + MAX\{100k_3, 150 + MAX(100k_1, 100k_2)\}$
 - AT at the output of the k₄-bit MUX and the third Carry-Out logic: $d_7 = MAX(d_4, d_6) + \varepsilon = 150 + MAX[100k_4, 150 + MAX\{100k_3, 150 + MAX(100k_1, 100k_2)\}]$

- 2) Compute the total delay when $k_1 = k_2 = k_3 = k_4$.
 - $k_1 = k_2 = k_3 = k_4 = 16.$
 - $d_7 = 150 + MAX[1600, 150 + MAX\{1600, 150 + MAX(1600, 1600)\}]$ = $150 + MAX[1600, 150 + MAX\{1600, 150 + 1600\}]$ = 150 + MAX[1600, 150 + 1750]= 150 + 1900 = 2050ps
- Compute k₁, k₂, k₃, and k₄ minimizing the delay and show the minimum delay. (Hint: (1) Use your intuition and some math. (2) If you want, you can program it to find k₁, k₂, k₃, k₄. In this case, you should show your program in your report).

1) We will first show that $k_1 = k_2$ will give us an optimal delay. Suppose MAX(100 k_1 , 100 k_2) affects the final delay. Then, $k_1 = k_2$ will minimize MAX(100 k_1 , 100 k_2), which will minimize the final delay.

 $k_1 = k_2$ leads to $d_7 = 150 + MAX[100k_4, 150 + MAX\{100k_3, 150 + 100k_1\}].$

By the same reason, $100k_3 = 150 + 100k_1$ will give us an optimal delay. From this, we get $k_3 = k_1 + 1.5$. However, k_3 should be an integer. If $k_3 \le k_1 + 1$, we get $d_7 = 150 + MAX[100k_4, 150 + 150 + 100k_1]$. However, reducing k_3 will increase k_4 , so let's set k_3 to $k_1 + 1$.

 $d_7 = 150 + MAX[100k_4, 300 + 100k_1]$. Setting $100k_4 = 300 + 100k_1$ will minimize $d_7 \rightarrow k_4 = 3 + k_1$

From $k_1 + k_2 + k_3 + k_4 = 64$, we get $k_1 + k_1 + (k_1 + 1) + (k_1 + 3) = 64 \rightarrow k_1 = 15$. $k_2 = 15$. $k_3 = 16$. $k_4 = 18$.

 $Delay = 150 + MAX[1800, 150 + MAX\{1600, 150 + MAX(1500, 1500)\}] = 1950 ps.$

2) We can also use a computer program to simulate this. The following C/C++ code simulates it.

```
#include <stdio.h>
int max (int a, int b) {
    if ( a > b )
        return a;
    return b;
}
```

```
int main () {
                                    int min_delay = 10000000; // min. delay achieved
                                    for (int k4 = 10; k4 \le 34; k4++) {
                                             for (int k3 = 10; k3 \le (44-k4); k3++) {
                                                      for (int k^2 = 10; k^2 \le (54 - k^4 - k^3); k^2 + +) {
                                                               int k1 = 64 - (k4 + k3 + k2);
                                                               int delay = 150 + \max(100 \times k4, 150 + \max(100 \times k3, 150 + \max(100 \times k1, 150 + \max(100 \times k1
 100*k2)));
                                                              if (delay <= min_delay) {
                                                                        printf ("(k4, k3, k2, k1, d) = (\%d, \%d, \%d, \%d, \%d)", k4, k3, k2,
k1, delay);
                                                                        min_delay = delay;
                                                                }
                                                       }
                                               }
                                     }
                                    return 0;
                            }
                           You can also download the following file:
                                    http://eecs.wsu.edu/~ee434/Homework/add.cpp
                           and compile it as follows (in the ee434-466 server):
                                    g++ add.cpp
                           which will generate a.out in your directory. Then, run it to see its usage:
                                    > ./a.out
                           The following shows the usage:
                                    ./a {delta_FA} {delta_MUX}
                           Run the program for the above problem as follows:
                                    ./a.out 100 150
                           which gives the following result (format: k4 k3 k2 k1 total delay in ps):
                                    delta_FA: 100 (ps) delta_MUX: 150 (ps)
                                    18 16 15 15 1950
                           You can also try some different combinations as follows:
                                    ./a.out 100 200
                           which gives the following result:
                                    17 17 15 15 2100
                                     18 16 15 15 2100
```

18 17 14 15	2100
18 17 15 14	2100
19 15 15 15	2100
19 16 14 15	2100
19 16 15 14	2100
19 17 13 15	2100
19 17 14 14	2100
19 17 15 13	2100

2. [Prefix Adder, **40 points**] Complete the following prefix adders by inserting merging blocks and drawing arrows (try to minimize the number of merging blocks inserted).



