

Computer Arithmetic Algorithms

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Number Representations

Information

- Textbook
 - Israel Koren, “Computer Arithmetic Algorithms,” Prentice Hall, 1993.

Outline

- Binary number system
- Radix conversion
- Negative numbers
 - Signed-magnitude
 - One's complement
 - Two's complement
- Addition and subtraction
- Arithmetic shift operations (sign extension)

The Binary Number System

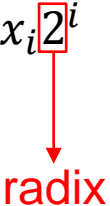
- A binary number of length n

$$(x_{n-1}, x_{n-2}, \dots, x_1, x_0)$$

- x_i : bit ($\in \{0,1\}$)

- An n -digit binary number

$$X = x_{n-1}2^{n-1} + x_{n-2}2^{n-2} + \dots + x_12^1 + x_02^0 = \sum_{i=0}^{n-1} x_i \boxed{2}^i$$


radix

- $x_i \in \{0,1\}$

- Example

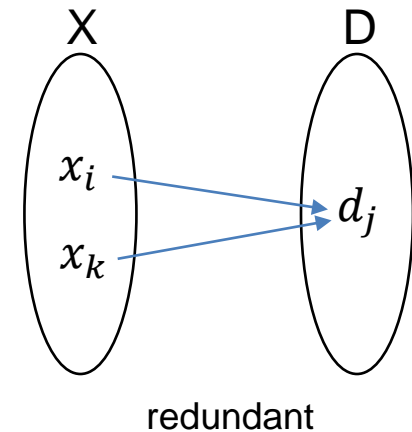
- $1011_2 = 1 \cdot 2^3 + 0 \cdot 2^2 + 1 \cdot 2^1 + 1 \cdot 2^0 = 11$

The Binary Number System

- Radix-10 numbers: decimal numbers
 - $(101)_{10} = 101$
- Radix-2 numbers: binary numbers
 - $(101)_2 = 5$
- Range of representable numbers
 - $[X_{min}, X_{max}]$
 - Example: The range of 4-bit unsigned binary numbers: $[0, 15]$
- Overflow
 - An arithmetic operation resulting in a number larger than X_{max} or smaller than X_{min}

The Binary Number System

- The conventional number systems
 - **Nonredundant**
 - Every number has a unique representation.
 - **Weighted**
 - $X = \sum_{i=0}^{n-1} x_i w_i$
 - w_i : weight
 - **Positional**
 - w_i depends only on the position of the digit x_i .
 - $w_i = r^i$
 - $0 \leq x_i \leq r - 1$
 - Otherwise (if $x_i \geq r$), the system becomes redundant.



The Binary Number System

- A mixed number (*fixed-point* representation)

$$\underbrace{(x_{k-1} x_{k-2} \dots x_1 x_0)}_{\text{integral part}} \cdot \underbrace{x_{-1} x_{-2} \dots x_{-m}}_{\text{fractional part}})_r$$

$$= (x_{k-1}r^{k-1} + x_{k-2}r^{k-2} + \dots + x_1r^1 + x_0r^0) + (x_{-1}r^{-1} + \dots + x_{-m}r^{-m})$$

$$= \sum_{i=-m}^{k-1} x_i r^i$$

- r^{-m} : indicates the position of the radix point.
- *ulp*: unit in the last position
 - $ulp = r^{-m}$

Radix Conversion

- Given

- $X = X_I + X_F = \sum_{i=0}^{k-1} x_i r^i + \sum_{i=-m}^{-1} x_i r^i$
 - X_I : integral part
 - X_F : fractional part
- r_D : destination number system

- Integral part

$$X_I = x_0 + r_D(x_1 + r_D(x_2 + \dots + r_D(x_{k-2} + r_D x_{k-1})))$$
$$0 \leq x_i < r_D$$

- $x_0 = X_I \text{ mod } r_D$
- $q_0 = x_1 + r_D(x_2 + r_D(x_3 + \dots))$
- $x_1 = q_0 \text{ mod } r_D$
- ...

Radix Conversion

- Fractional part

$$X_F = r_D^{-1}(x_{-1} + r_D^{-1}(x_{-2} + \cdots + r_D^{-1}(x_{-(m-1)} + r_D^{-1}x_{-m})))$$
$$0 \leq x_i < r_D$$

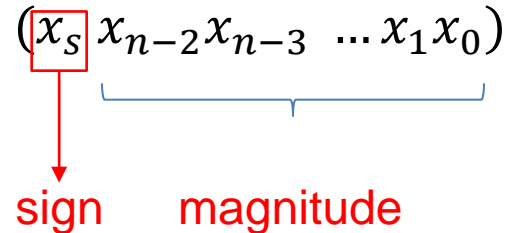
- x_{-1} : the integral part of $X_F r_D$
- $m_{-1} = r_D^{-1}(x_{-2} + \cdots + r_D^{-1}(x_{-(m-1)} + r_D^{-1}x_{-m}))$
- x_{-2} : the integral part of $m_{-1} r_D$
- ...

- Example

- $46.375_{10} = 101110.011_2$

Negative Numbers: Signed-Magnitude

- Representation



- Positive numbers: $0xxxxx$
 - $0xxxxx_2$
- Negative numbers: $(r-1)xxxxx$
 - $1xxxxx_2$
- Range of positive numbers for $(s x_{k-2} x_{k-3} \dots x_1 x_0 . x_{-1} x_{-2} \dots x_{-m})$
 - $[0, r^{k-1} - ulp]$
- Range of negative numbers for $(s x_{k-2} x_{k-3} \dots x_1 x_0 . x_{-1} x_{-2} \dots x_{-m})$
 - $[-(r^{k-1} - ulp), 0]$

Negative Numbers: Signed-Magnitude

- Problems
 - Two zeros
 - 0000...0
 - 1000...0
 - Difficulties in subtraction
 - If $X > Y$, $X - Y$ is ok.
 - However, $Y - X$ is $-(X - Y)$, so we should compute $X - Y$, then set the MSB to 1.

Negative Numbers: Complement

- Represent $-Y$ by $R - Y$.
- Then, $-(-Y) = R - (R - Y) = Y$.
- If $X > Y > 0$, $Y - X = Y + (R - X) = R - (X - Y) = X - Y$.
- Two requirements
 - When $X > Y > 0$, $X - Y = X + (R - Y) = R + (X - Y)$ should be $X - Y$.
 - Computation of $R - Y$ should be simple.
- Notation

$$\begin{aligned}\bar{x}_i &= (r - 1) - x_i \\ \bar{X} &= (\overline{x_{k-1}}, \overline{x_{k-2}}, \dots, \overline{x_0}, \overline{x_{-1}}, \dots, \overline{x_{-m}}) \\ X + \bar{X} &= ((r - 1), (r - 1), \dots, (r - 1)) \\ X + \bar{X} + ulp &= r^k\end{aligned}$$

Negative Numbers: Complement

- One's complement (Diminished-radix complement)
 - $R = r^k - ulp$
 - $-X = R - X = \bar{X}$
 - When $X > Y > 0$, $Y - X = Y + (R - X) = R - (X - Y) = -(Y - X)$
 - When $X > Y > 0$, $X - Y = X + (R - Y) = R + (X - Y) = r^k + (X - Y) - ulp = (X - Y) - ulp$
 - This requires a correction step (add ulp).
 - When $X > 0$, $Y > 0$, $-X - Y = (R - X) + (R - Y) = R + R - (X + Y) = r^k + R - (X + Y) = R - (X + Y) = -(X + Y)$
- Two's complement (Radix complement)
 - $R = r^k$
 - $-X = R - X = \bar{X} + ulp$
 - When $X > Y > 0$, $Y - X = Y + (R - X) = R - (Y - X) = -(Y - X)$
 - When $X > Y > 0$, $X - Y = X + (R - Y) = R + (X - Y) = r^k + (X - Y) = X - Y$
 - When $X > 0$, $Y > 0$, $-X - Y = (R - X) + (R - Y) = R + R - (X + Y) = r^k + R - (X + Y) = R - (X + Y) = -(X + Y)$

Two's Complement Representation

- One zero
 - 000...0
- A negative number that does not have a positive equivalent
 - 100...0
- Range
 - $[-2^{n-1}, 2^{n-1} - ulp]$

- $X = -x_{n-1} \cdot 2^{n-1} + \sum_{i=0}^{n-2} x_i 2^i$

Proof) If $x_{n-1} = 0$ (positive), $X = \sum_{i=0}^{n-2} x_i 2^i = -x_{n-1} \cdot 2^{n-1} + \sum_{i=0}^{n-2} x_i 2^i$.

If $x_{n-1} = 1$ (negative), $X = -[\bar{X} + ulp] = -[\sum_{i=0}^{n-1} \bar{x}_i 2^i + 1]$

$$= -[\sum_{i=0}^{n-1} (1 - x_i) 2^i + 1] = -[2^n - 1 - \sum_{i=0}^{n-1} x_i 2^i + 1]$$

$$= -2^n + \sum_{i=0}^{n-1} x_i 2^i = -2^{n-1} + \sum_{i=0}^{n-2} x_i 2^i$$

One's Complement Representation

- Two zeros

- 000...0
- 111...1

- Range

- $[-(2^{n-1} - ulp), (2^{n-1} - ulp)]$

- $X = -x_{n-1}(2^{n-1} - ulp) + \sum_{i=0}^{n-2} x_i 2^i$

Proof) If $x_{n-1} = 0$ (positive), $X = \sum_{i=0}^{n-2} x_i 2^i = -x_{n-1} \cdot (2^{n-1} - ulp) + \sum_{i=0}^{n-2} x_i 2^i$.

$$\begin{aligned} \text{If } x_{n-1} = 1 \text{ (negative), } X &= -\bar{X} = -\sum_{i=0}^{n-1} \bar{x}_i 2^i = -\sum_{i=0}^{n-1} (1 - x_i) 2^i \\ &= -[2^n - 1 - \sum_{i=0}^{n-1} x_i 2^i] = -2^n + 1 + \sum_{i=0}^{n-1} x_i 2^i \\ &= -2^{n-1} + 1 + \sum_{i=0}^{n-2} x_i 2^i = -x_{n-1} \cdot (2^{n-1} - ulp) + \sum_{i=0}^{n-2} x_i 2^i \end{aligned}$$

Addition and Subtraction

- Signed-magnitude

$$\begin{array}{r} 0\ 1\ 0\ 1\ 1\ (11) \\ +\ 0\ 0\ 0\ 1\ 0\ (2) \\ \hline =\ 0\ 1\ 1\ 0\ 1\ (13) \end{array}$$

$$\begin{array}{r} \boxed{0}\ 1\ 0\ 1\ 1\ (11) \\ +\ 0\ 0\ 1\ 1\ 0\ (6) \\ \hline =\ \boxed{1}\ 0\ 0\ 0\ 1\ (-1) \end{array}$$

↓
overflow

$$\begin{array}{r} \boxed{1}\ 0\ 0\ 1\ 1\ (-3) \\ +\ 1\ 0\ 1\ 0\ 1\ (-5) \\ \hline =\ 1\ \boxed{0}\ 1\ 0\ 0\ 0\ (+8) \end{array}$$

↓
overflow

$$\begin{array}{r} 1\ 1\ 0\ 0\ 0\ (-8) \\ +\ 1\ 1\ 0\ 0\ 1\ (-9) \\ \hline =\ \boxed{1\ 1}\ 0\ 0\ 0\ 1\ (-1) \end{array}$$

↓
overflow

$$\begin{array}{r} 0\ 1\ 1\ 0\ 1\ (13) \\ +\ 1\ 1\ 0\ 0\ 0\ (-8) \\ \hline =\ 1\ 0\ 0\ 1\ 0\ 1\ (5) \end{array}$$

$$\begin{array}{r} 1\ 1\ 1\ 0\ 1\ (-13) \\ +\ 0\ 1\ 0\ 0\ 0\ (8) \\ \hline =\ 1\ 0\ 0\ 1\ 0\ 1\ (+5) \end{array}$$

- As shown above, addition and subtraction of signed-magnitude numbers need corrections.

Addition and Subtraction

- Two's complement

$$\begin{array}{r} 01011 \text{ (11)} \\ + 00010 \text{ (2)} \\ \hline = 01101 \text{ (13)} \end{array}$$

$$\begin{array}{r} 01011 \text{ (11)} \\ + 00110 \text{ (6)} \\ \hline = 10001 \text{ (-15)} \end{array}$$

↓
overflow
indicator

$$\begin{array}{r} 11101 \text{ (-3)} \\ + 11011 \text{ (-5)} \\ \hline = 111000 \text{ (-8)} \end{array}$$

$$\begin{array}{r} 11000 \text{ (-8)} \\ + 10111 \text{ (-9)} \\ \hline = 101111 \text{ (-17)} \end{array}$$

↓
overflow
indicator

$$\begin{array}{r} 01101 \text{ (13)} \\ + 11000 \text{ (-8)} \\ \hline = 100101 \text{ (5)} \end{array}$$

$$\begin{array}{r} 10011 \text{ (-13)} \\ + 01000 \text{ (8)} \\ \hline = 11011 \text{ (-5)} \end{array}$$

Arithmetic Shift Operations

- For given $\{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}$

- Signed-magnitude

$$x_{n-1}, 0, 0, \dots, 0, \{x_{n-2}, \dots, x_1, x_0\}, 0, 0, \dots$$

- Two's complement

$$x_{n-1}, x_{n-1}, \dots, x_{n-1}, \{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}, 0, 0, \dots$$

- One's complement

$$x_{n-1}, x_{n-1}, \dots, x_{n-1}, \{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}, x_{n-1}, x_{n-1}, \dots$$

Supplementary Materials

Arithmetic Shift Operations

- Derivation of the extension of the two's complement

$$X = X_{n-1\dots 0} = \{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}$$

If $X > 0$, it is the same as the signed-magnitude case.

If $X < 0$ (if $X \neq -2^n$),

$$X = X_{n-1\dots 0} = \{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}$$

$$-X = \{\overline{X_{n-1\dots 0}} + ulp_0\}$$

$$(-X)_{\text{ext}} = \overline{00 \dots 0\{\overline{X_{n-1\dots 0}} + ulp_0\}00 \dots 0} = 00 \dots 0\{K\}00 \dots 0$$

$$-(-X)_{\text{ext}} = \overline{00 \dots 0\{K\}00 \dots 0} + ulp_{-m} = 11 \dots 1\{\overline{K}\}11 \dots 1 + ulp_{-m}$$

$$= 11 \dots 1\{\overline{K} + ulp_0\}00 \dots 0 = 11 \dots 1\{2^n - K\}00 \dots 0$$

$$= 11 \dots 1\{2^n - (\overline{X_{n-1\dots 0}} + ulp_0)\}00 \dots 0 = 11 \dots 1\{X_{n-1\dots 0}\}00 \dots 0$$

$$\therefore X_{\text{ext}} = x_{n-1} \dots x_{n-1}\{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}00 \dots 0$$

If $X = 2^n$, we can directly prove it.

Arithmetic Shift Operations

- Derivation of the extension of the one's complement

$$X = X_{n-1\dots 0} = \{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}$$

If $X > 0$, it is the same as the signed-magnitude case.

If $X < 0$,

$$X = X_{n-1\dots 0} = \{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}$$

$$-X = \{\overline{X_{n-1\dots 0}}\}$$

$$(-X)_{\text{ext}} = \overline{00 \dots 0 \{X_{n-1\dots 0}\} 00 \dots 0}$$

$$-(-X)_{\text{ext}} = \overline{00 \dots 0 \{\overline{X_{n-1\dots 0}}\} 00 \dots 0} = 11 \dots 1 \{X_{n-1\dots 0}\} 11 \dots 1$$

$$\therefore X_{\text{ext}} = x_{n-1} \dots x_{n-1} \{x_{n-1}, x_{n-2}, \dots, x_1, x_0\} x_{n-1} \dots x_{n-1}$$