

EE 466/586 VLSI Design

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Two-Phase Systems Lecture 20

Clocking disciplines

- Rules for constructing sequential machines.
 - Combinations of registers and gates.
 - Behavior of clocks and primary inputs over time.
- Rules are sufficient to guarantee that the system will work at some clock rate.
 - May not be as fast as we want.

Clocking Rules

- Combinational logic gates cannot be connected in a cycle
- All components must have bounded delay

Flip-flop rules

- Primary inputs change after clock (\$) edge.
- Primary inputs must stabilize before next clock edge.
- Rules allow changes to propagate through combinational logic for next cycle.
- Flip-flop outputs hold current-state values for next-state computation.



Latch-based machines

- Latches do not cut combinational logic when clock is active.
- Latch-based machines must use multiple ranks of latches.
- Multiple ranks require multiple phases of clock.

Two-sided latch constraint

- Latch must be open less than the shortest combinational delay.
- Period between latching operations must be longer than the longest combinational delay.
- Note: difference between shortest and longest combinational delay may be large (sum₀ vs. sum₃₁).

Strict two-phase clocking discipline

- Strict two-phase discipline is conservative but works.
- Can be relaxed later with proper knowledge of constraints.
- Strict two-phase machine makes latch-based machine behave more like flip-flop design, but requires multiple phases.





Why it works

- Each phase has a one-sided constraint: phase must be long enough for all combinational delays.
- If there are no combinational loops, phases can always be stretched to make that section of the machine work.
- Total clock period depends on sum of phase periods.

Clocking types

- Logic on different phases operate at different times—can't mix signals from different phases.
- Primary inputs must obey the same rules as internal signals.
- Clocking types help us to ensure that machine structure is valid.

Stable signals

- A logic signal is always stable during one phase—phase in which the latch which produced it is not active.
- Easiest to think of machine behavior in terms of stable signals, though signals propagate while not stable.

Signal types

- Clocks are separate type: ϕ_1 , ϕ_2 .
- Two types of stable data signal:
 - stable ϕ_1 (s ϕ_1)
 - stable ϕ_2 (s ϕ_2)
- A stable signal has a complementary valid signal:

- stable ϕ_2 (s ϕ_2) = valid ϕ_1 (v ϕ_1)



How clocking types combine



Clocking types in the two-phase machine $I_1(s \phi_2)$ **s** φ₂ combinational logic D Q $O_1(s \phi_2)$



Example: shift register

- Want to displace bit by n registers in n cycles.
- Each register requires two phases:



Shift register operation



 $\phi_1 = 1, \phi_2 = 0$



 $\phi_1 = 0, \, \phi_2 = 1$

Non-strict disciplines

Some relaxation of the rules can be useful:

- reduce area;
- increase performance.
- Rules must be relaxed in a way that ensures the machine will still work.

Use logic to generate a clock signal which is not always active.

Qualified clocks

- Qualification must not introduce glitches into the clock—glitches violate the fundamental definition of a clock by introducing extra edges.
- Use stable signals to qualify clocks.



Clock logically combined with signal:



Uses of qualified clocks

- May want to conditionally load a register.
- May qualify a clock to turn off machine for low-power operation.
- Latch must not lose its value during inactive period.
- Difficult to ensure that logic value will come high in time—use quasi-static latch.

Qualified clocks and skew

- Logic in the clocking path introduces delay.
- Delay can cause clock to arrive at latches at different times, violating clocking assumptions.
- When designing qualification logic:
 - minimize and check skew;
 - sharpen clock edge.

Unbalanced delays

Logic with unbalanced delays leads to inefficient use of logic:



Flip-flop-based system performance analysis D Q clock combinational logic

Flip-flop-based system model

- Clock signal is perfect (no rise/fall), period P.
- Clock event on rising edge.
- Setup time s.
 - Time from arrival of combinational logic event to clock event.
- Propagation time p.
 - Time for value to go from flip-flop input to output.
- Worst-case combinational delay C.
 - Time from output of flip-flop to input.





Clock period constraint ■ P >= C + s + p. p C S D Q clock combinational logic

Clock with rise/fall





Latch system clock period

- For each phase, phase period must be longer than sum of:
 - combinational delay;
 - latch propagation delay.
- Phase period depends on longest path.



Two-phase timing parameters



Clock period constraint

• Total clock period (both phases): $-P \ge C1 + C2 + 2s + 2p.$

Each phase must meet timing for its own latch.

Latch-based system model D Q D Q φ₂

Example with unbalanced phases

One phase is much longer than the other:





Compute only part of long paths in one phase:



Spreading out a phase, cont'd.

Use other phase for end of long logic block and all of short logic block:





- **Skew**: relative delay between events.
- Signal skew: most important for asynchronous, timing-dependent logic.
- Clock skew: can harm any sequential system.



Machine data signals must obey setup and hold times—avoid signal skew.





Clock must arrive at all memory elements in time to load data.



Clock skew example





Clock distribution

- Often one of the hardest problems in clock design.
 - Fast edges.
 - Minimum skew.

