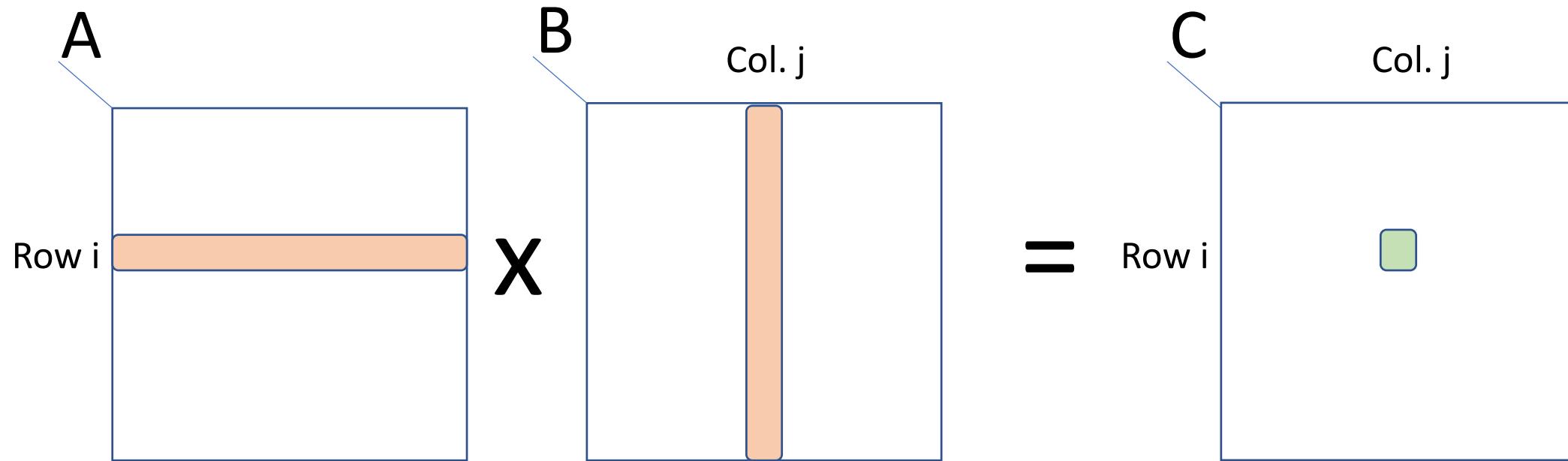


# Cannon's Algorithm for Matrix Multiplication

L. Cannon, 1969

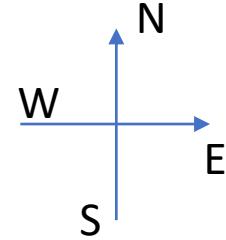
For ease of exposition, let us assume square matrices/

$$\text{Input: } A(n \times n) \times B(n \times n) = C(n \times n)$$



$$C[i][j] = \sum_{k=0}^{n-1} A[i][k] \times B[k][j]$$

# Cannon's algorithm



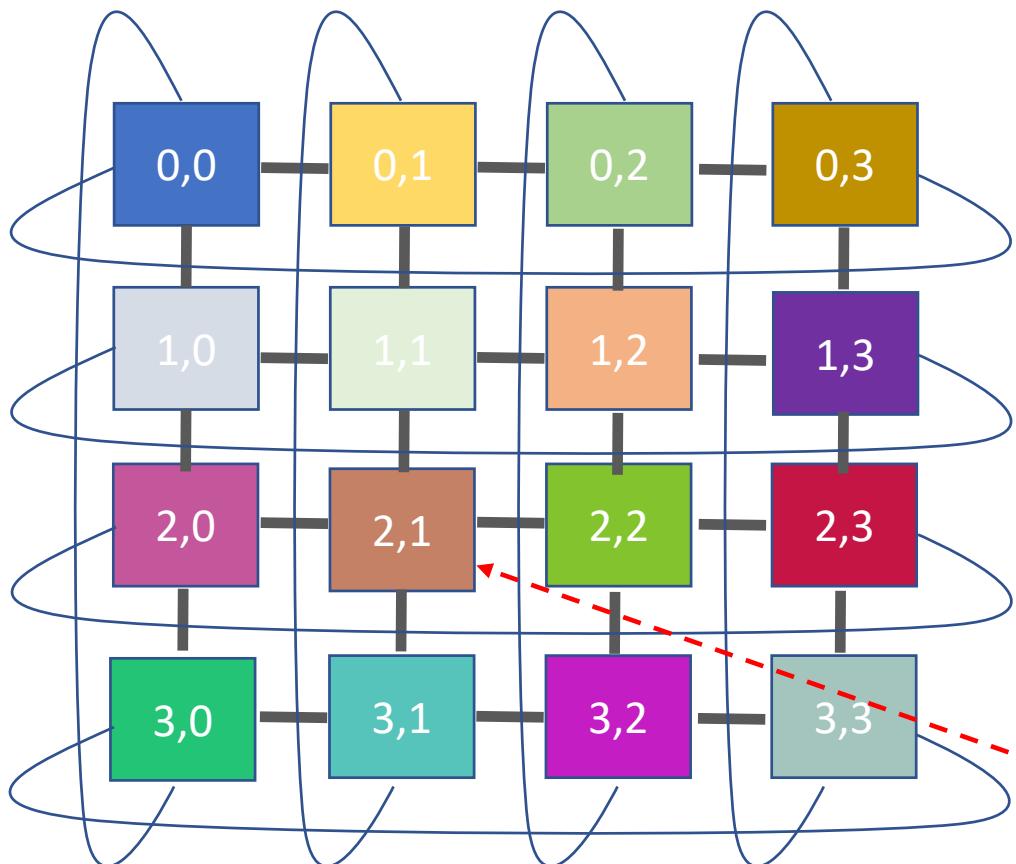
## // Initialize matrices A and B on the processes:

- First load cells  $A[i,j]$  and  $B[i,j]$  on proc.  $[i,j]$
- Circular left shift the  $i^{\text{th}}$  row of matrix A,  $i$  times
- Circular up shift the  $j^{\text{th}}$  column of matrix B,  $j$  times
- Set **a** and **b** to the corresponding values of A and B the proc. got after the shift.

## // Main algorithm @ proc. $[i,j]$

- $c=0$  // local **c** element which will be output by this proc.
- For the length of a row:
  - $c += a \times b$
  - Move **a** element one step left (west)
  - Move **b** element one step up (north)
- Output **c** // this will be the output cell  $C[i,j]$

# Assume that the Processes are Logically Arranged as a Torus (wrap-around mesh)



$$P = 16 = 4 \times 4$$

- ❖ Let us assume we have  $n^2=p$
  - ❖ Identify processes by their tile coordinate:
    - E.g.,  $(i,j)$  is the label of the process in row  $i$  and column  $j$  on the Torus
- 

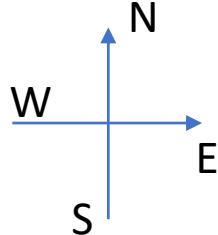
## Strategy:

- $C[i][j]$  will be generated by process  $(i,j)$

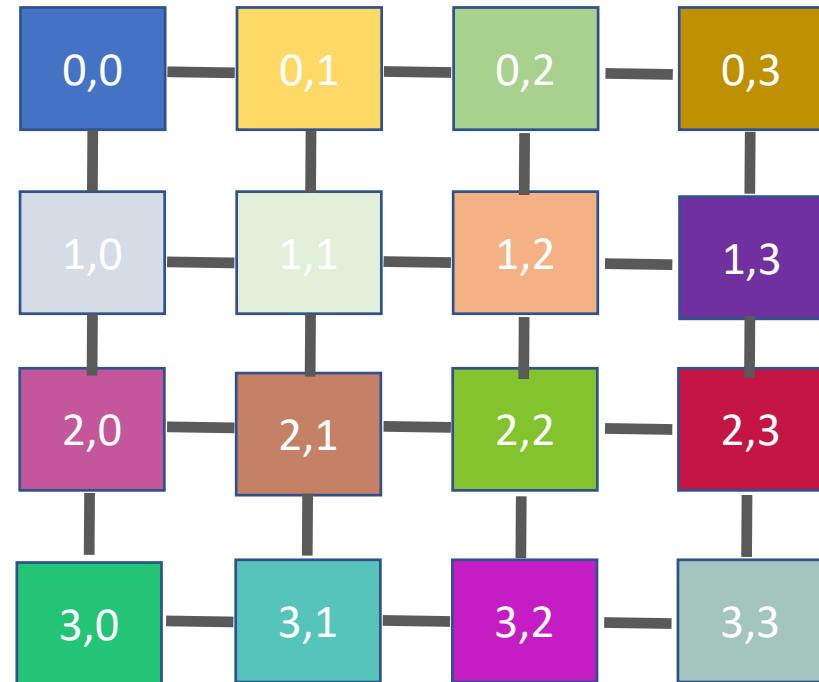
e.g.,

Proc. (2,1) will generate output value  $C[2][1]$

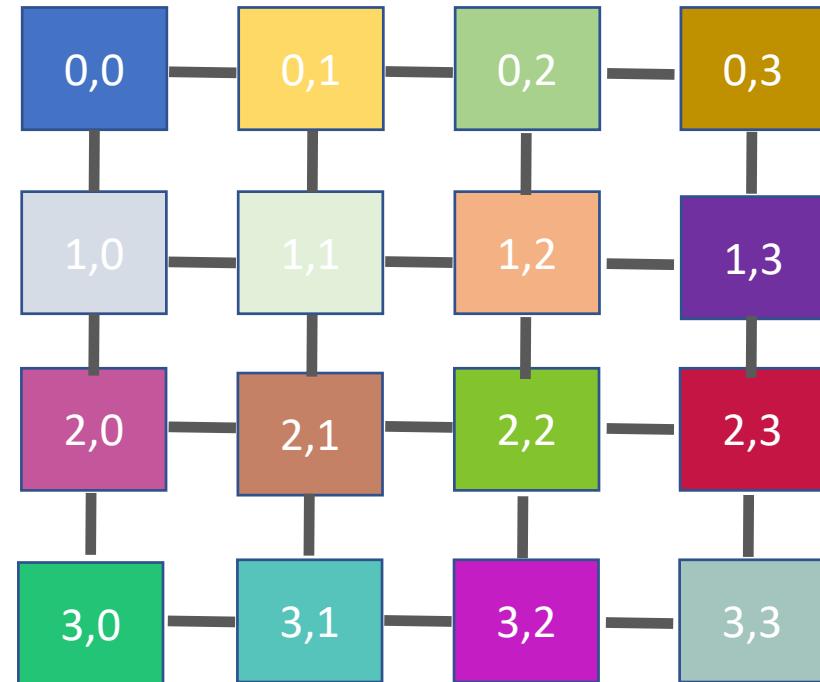
Load the matrices A and B are loaded such that cell  $(i,j)$  resides on proc.  $(i,j)$



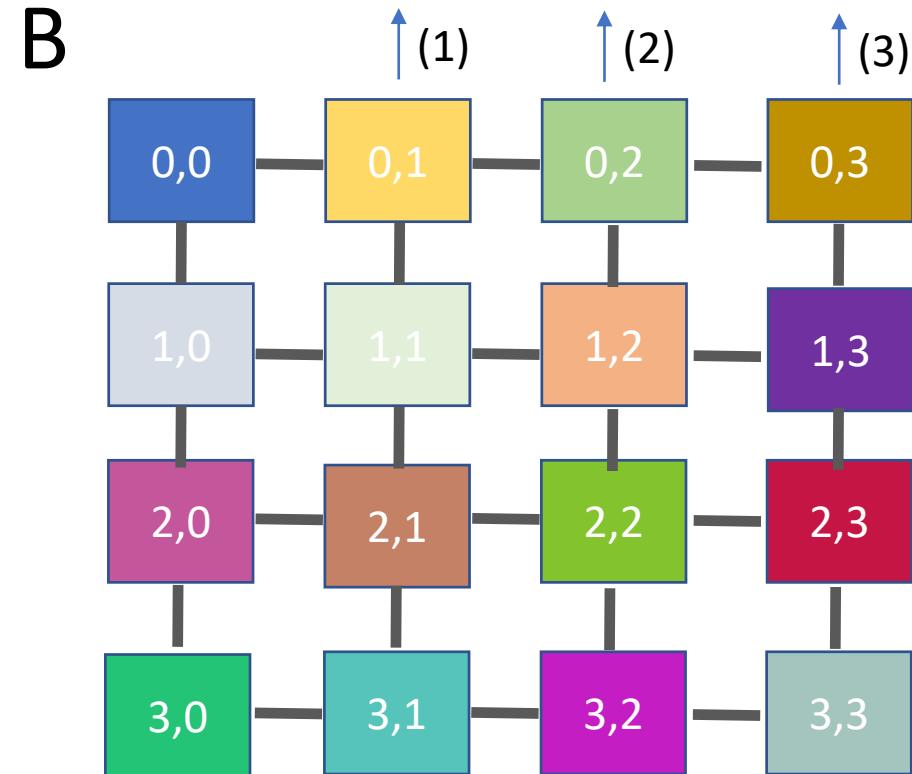
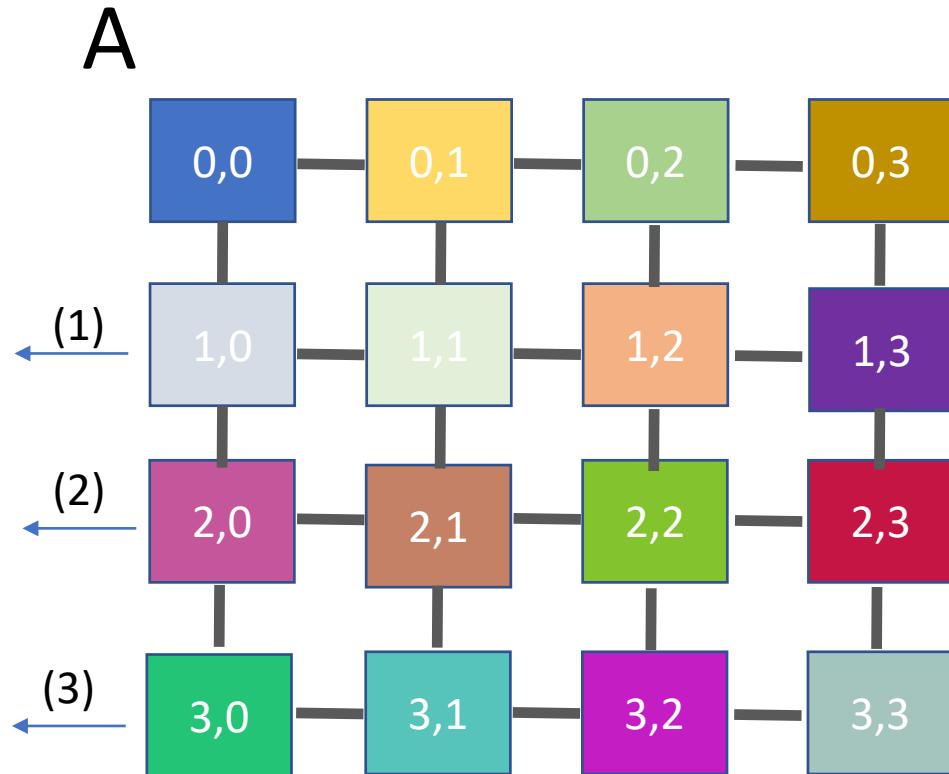
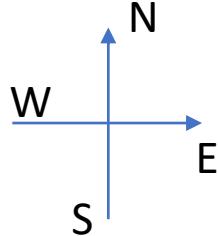
A



B

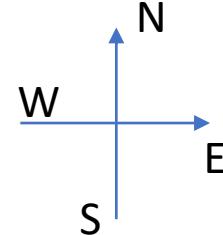


Move row  $i$  of A,  $i$  times West (circular left shift)  
Move col.  $j$  of B,  $j$  times North (circular up shift)

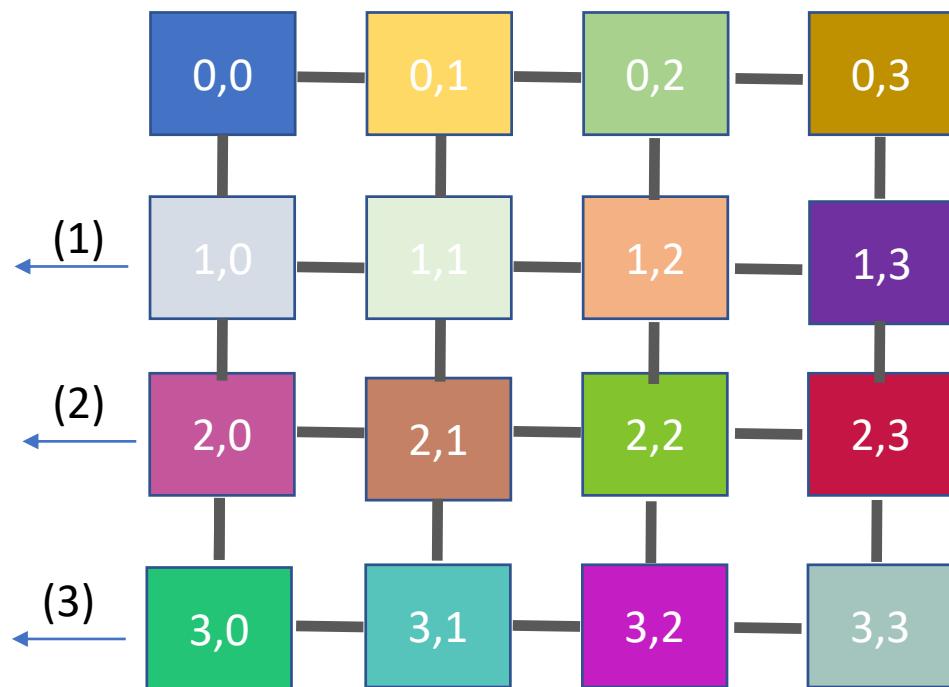


Move row  $i$  of A,  $i$  times West (circular left shift)

Move col.  $j$  of B,  $j$  times North (circular up shift)

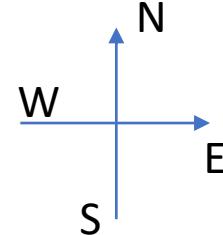


A

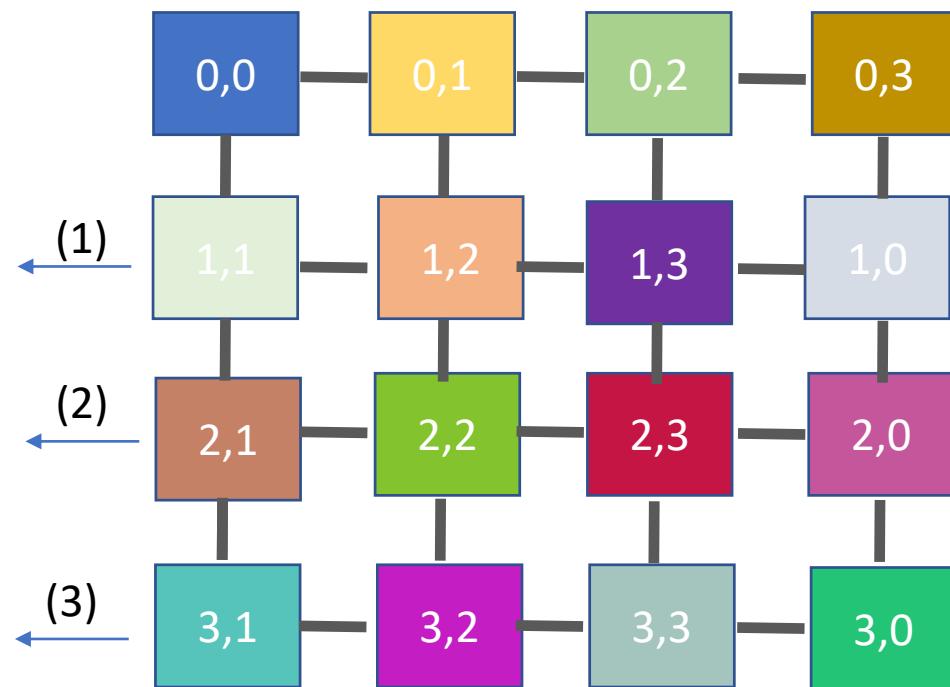


Move row  $i$  of  $A$ ,  $i$  times West (circular left shift)

Move col.  $j$  of  $B$ ,  $j$  times North (circular up shift)

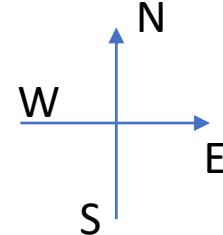


$A$

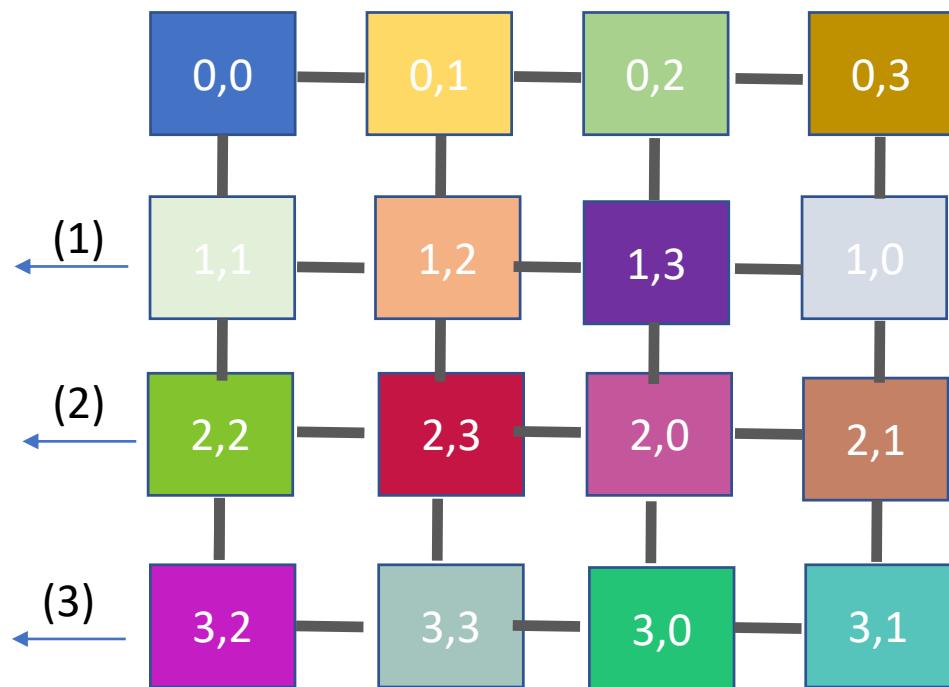


Move row  $i$  of  $A$ ,  $i$  times West (circular left shift)

Move col.  $j$  of  $B$ ,  $j$  times North (circular up shift)

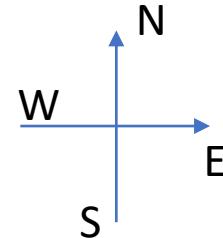


$A$

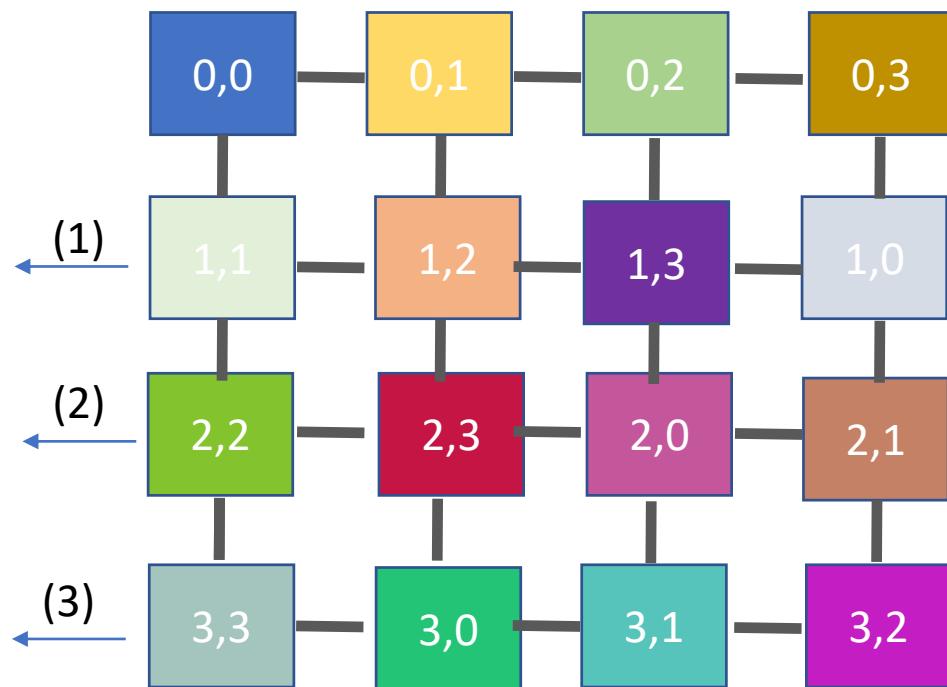


Move row  $i$  of  $A$ ,  $i$  times West (circular left shift)

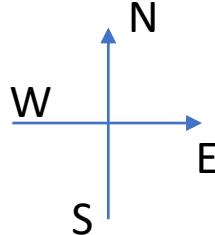
Move col.  $j$  of  $B$ ,  $j$  times North (circular up shift)



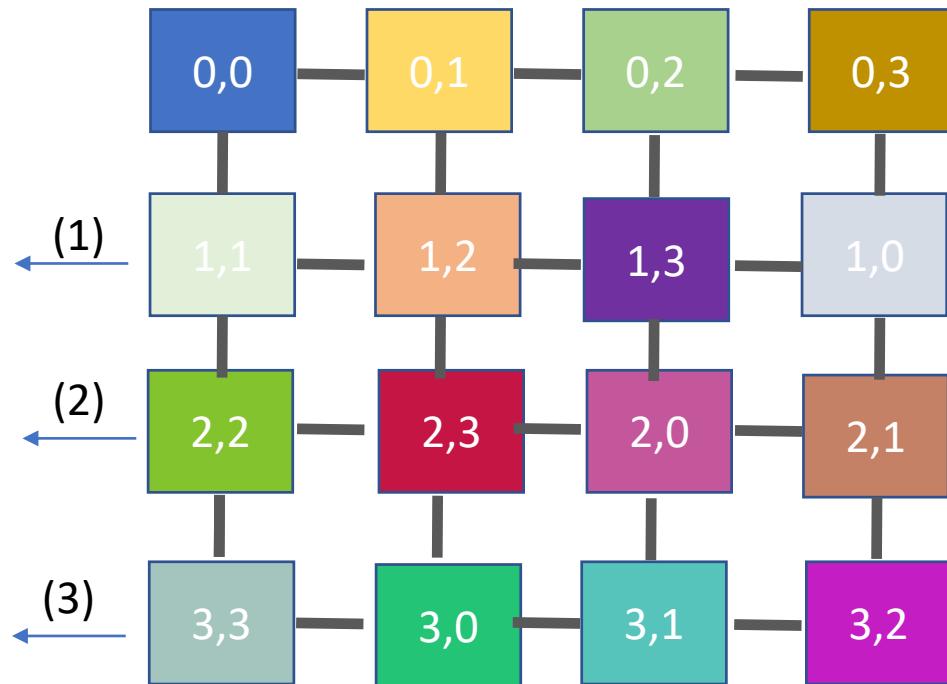
$A$



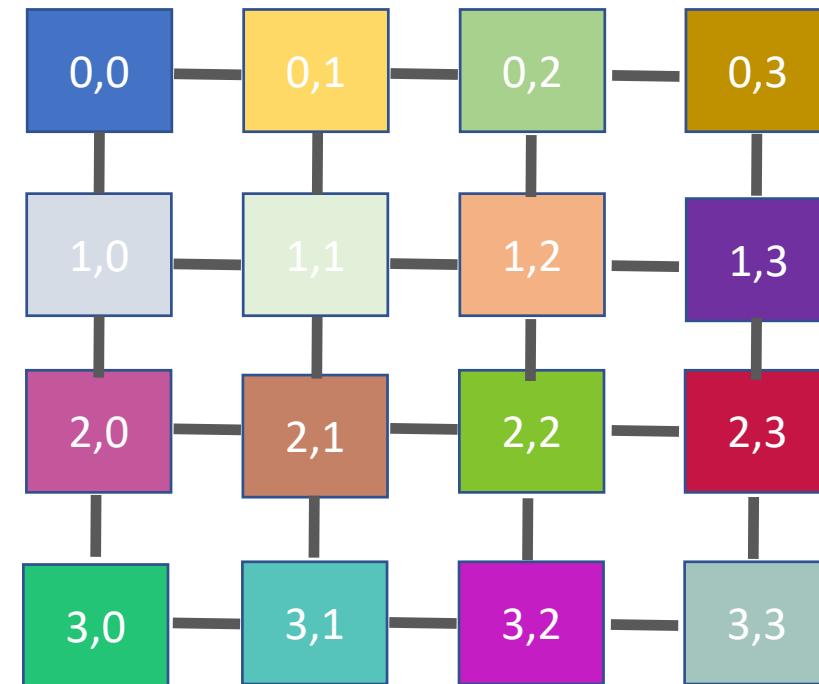
Move row  $i$  of A,  $i$  times West (circular left shift)  
Move col.  $j$  of B,  $j$  times North (circular up shift)



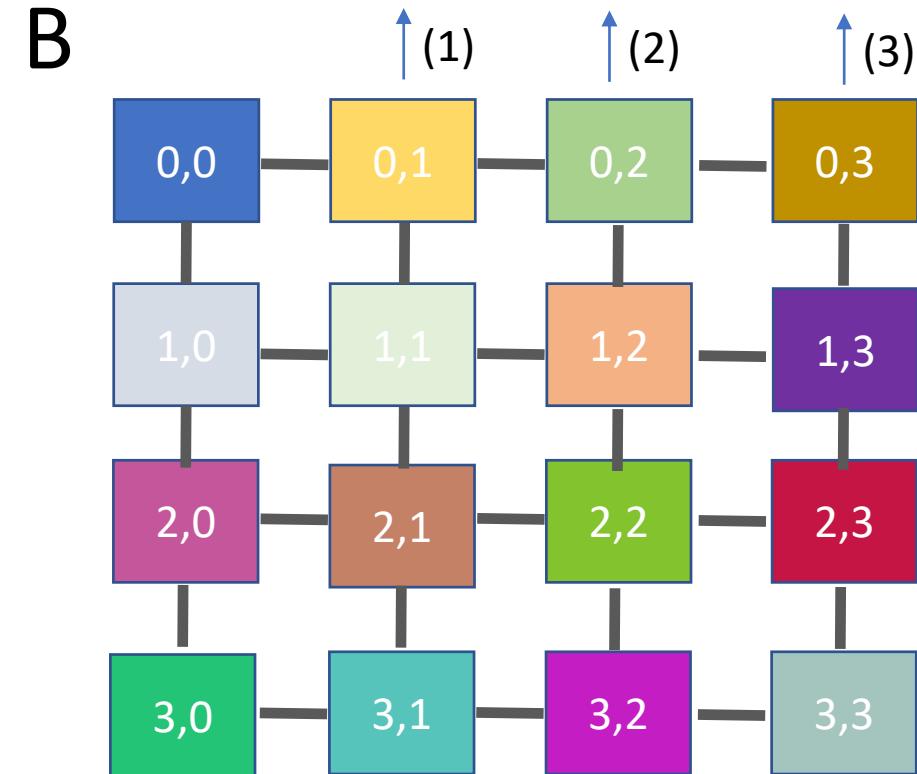
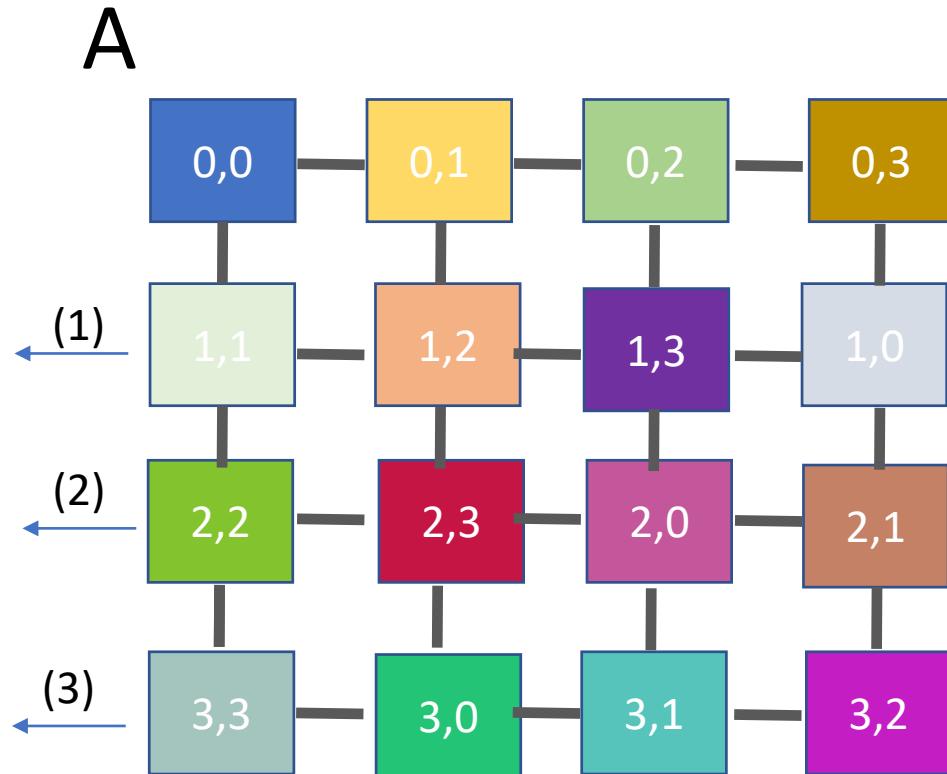
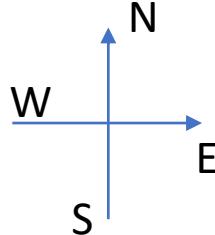
A



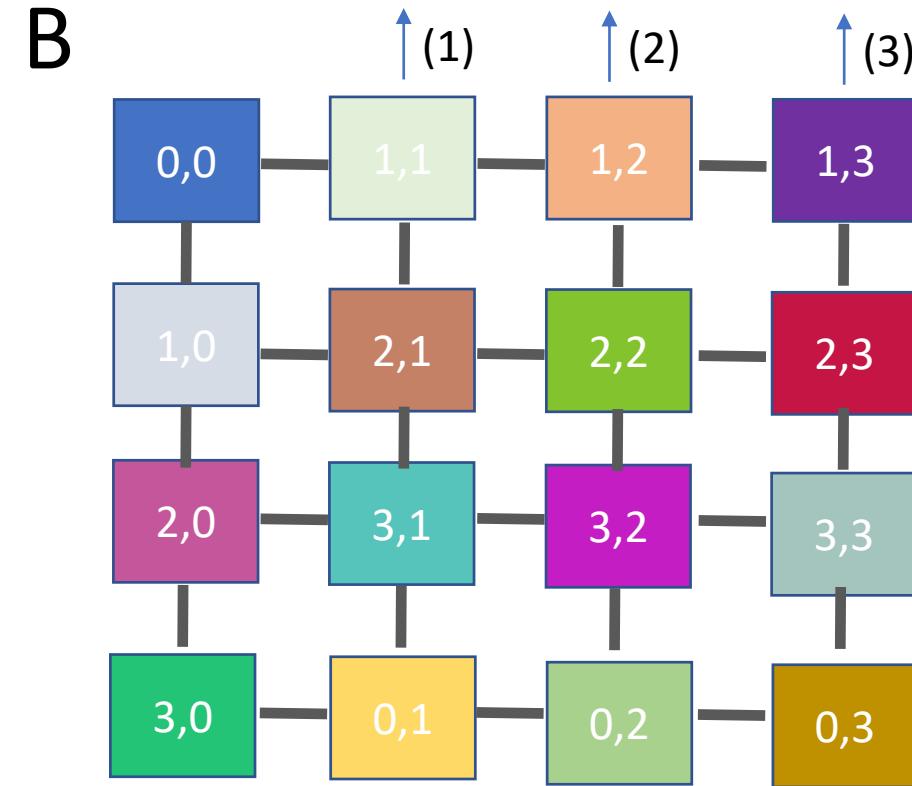
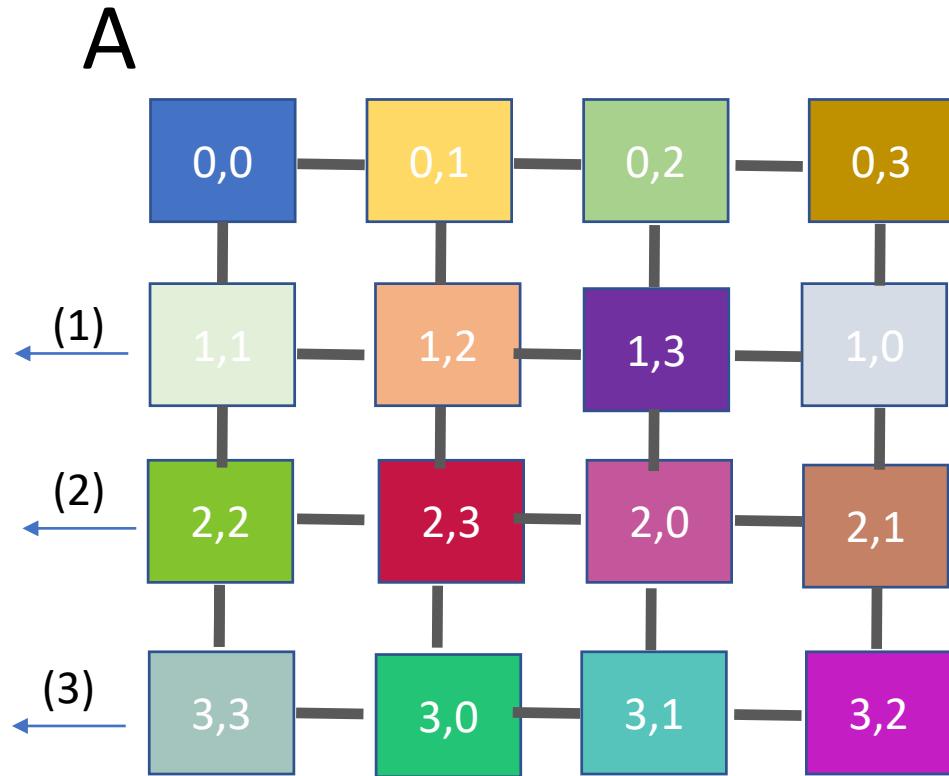
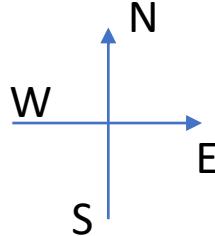
B



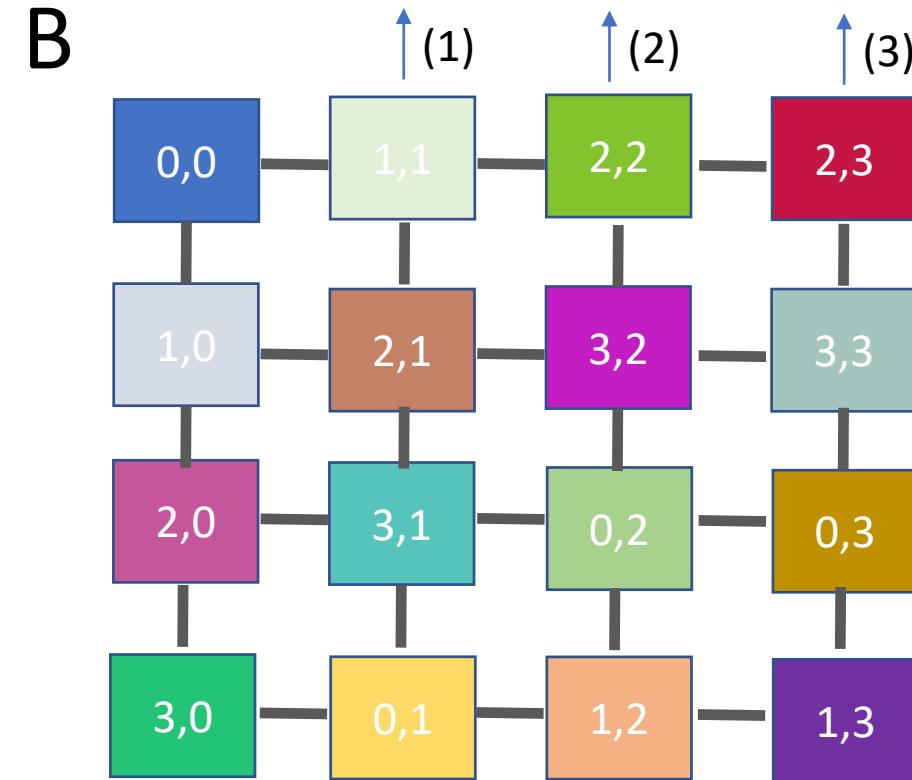
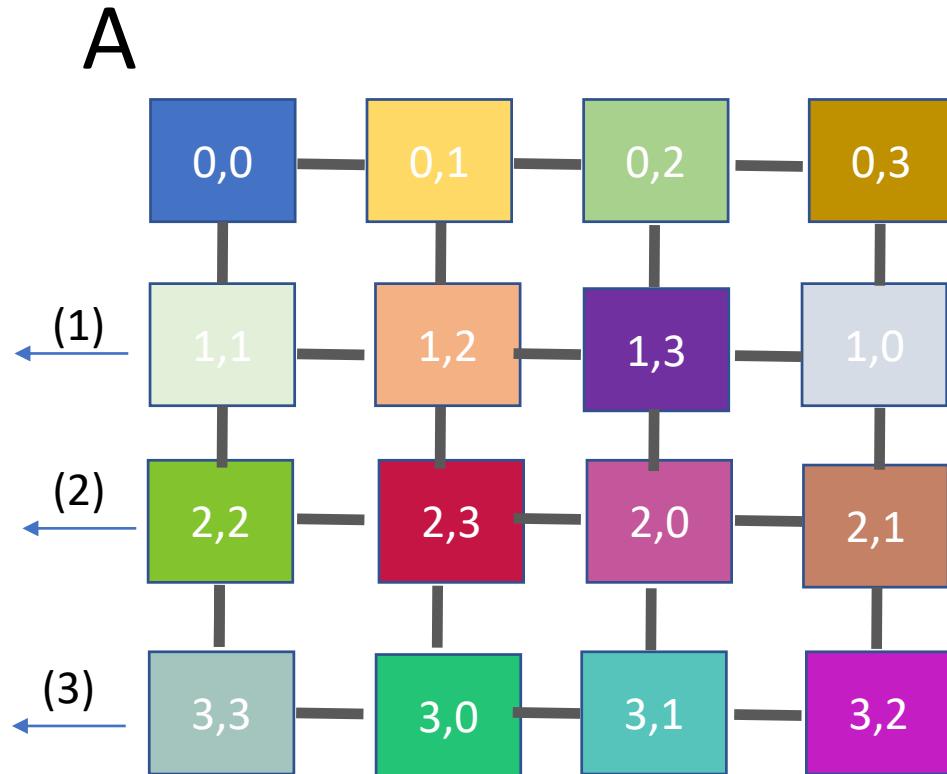
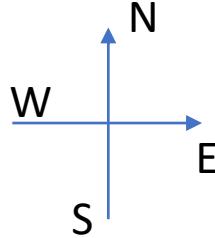
Move row  $i$  of A,  $i$  times West (circular left shift)  
Move col.  $j$  of B,  $j$  times North (circular up shift)



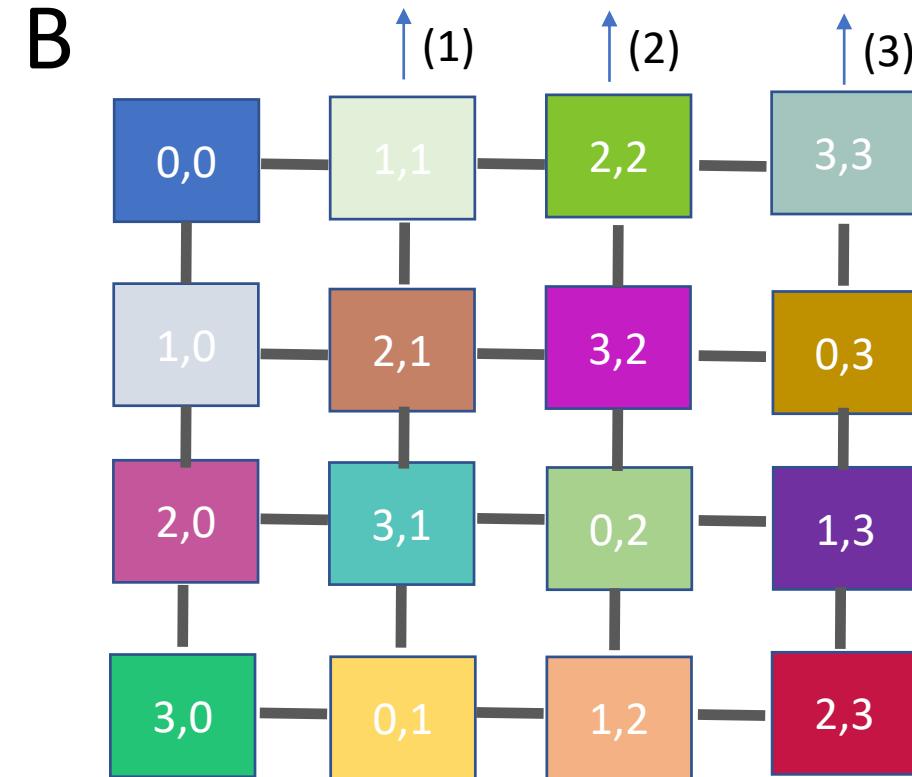
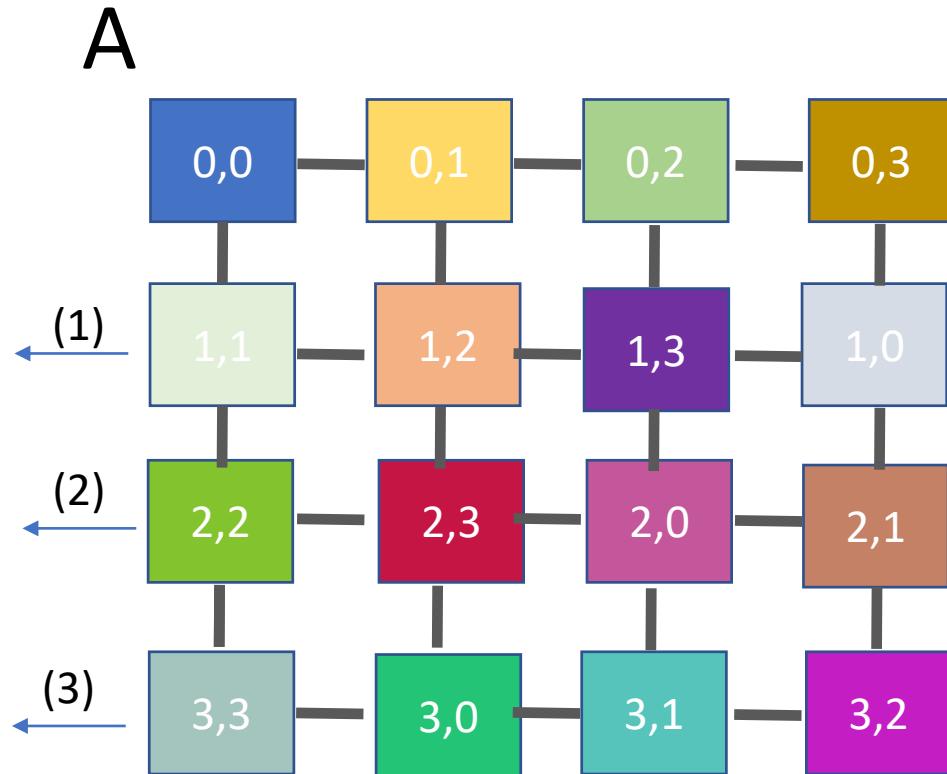
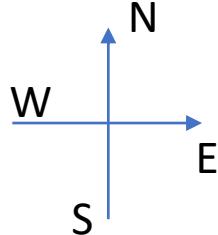
Move row  $i$  of A,  $i$  times West (circular left shift)  
Move col.  $j$  of B,  $j$  times North (circular up shift)



Move row  $i$  of A,  $i$  times West (circular left shift)  
Move col.  $j$  of B,  $j$  times North (circular up shift)

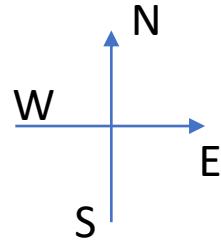


Move row  $i$  of A,  $i$  times West (circular left shift)  
Move col.  $j$  of B,  $j$  times North (circular up shift)

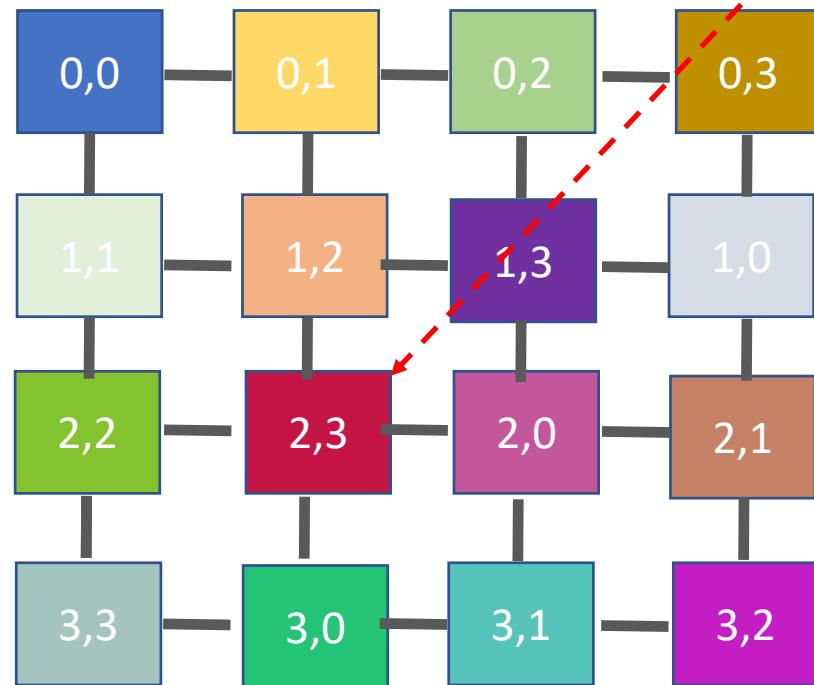


# Main Algorithm

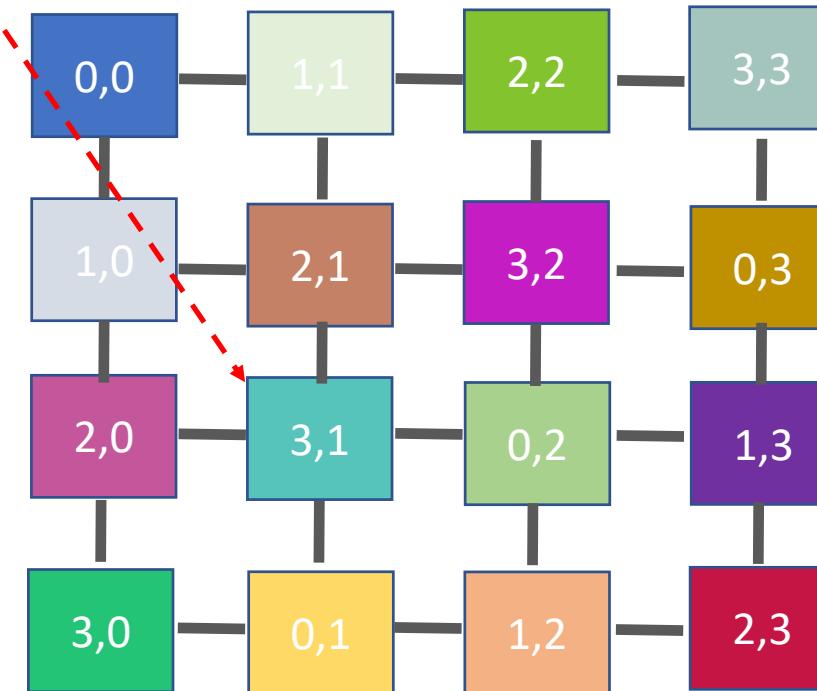
Let us focus on the computation of  $C[2][1]$ .  
Tile (2,1) will compute it.



A

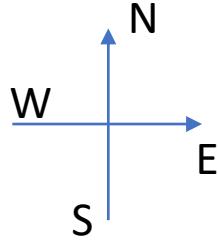


B

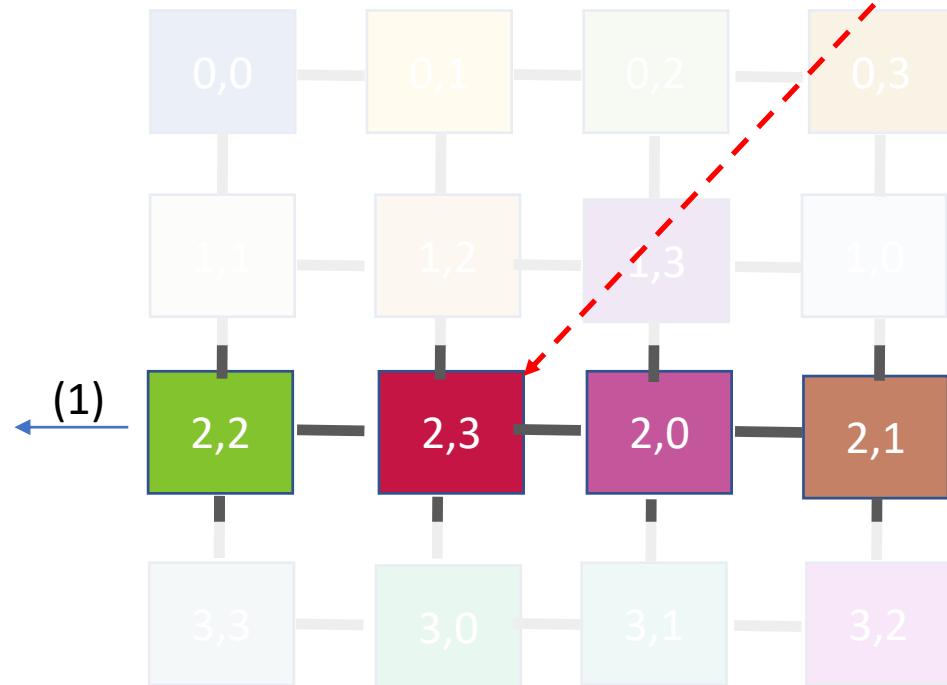


# Main Algorithm

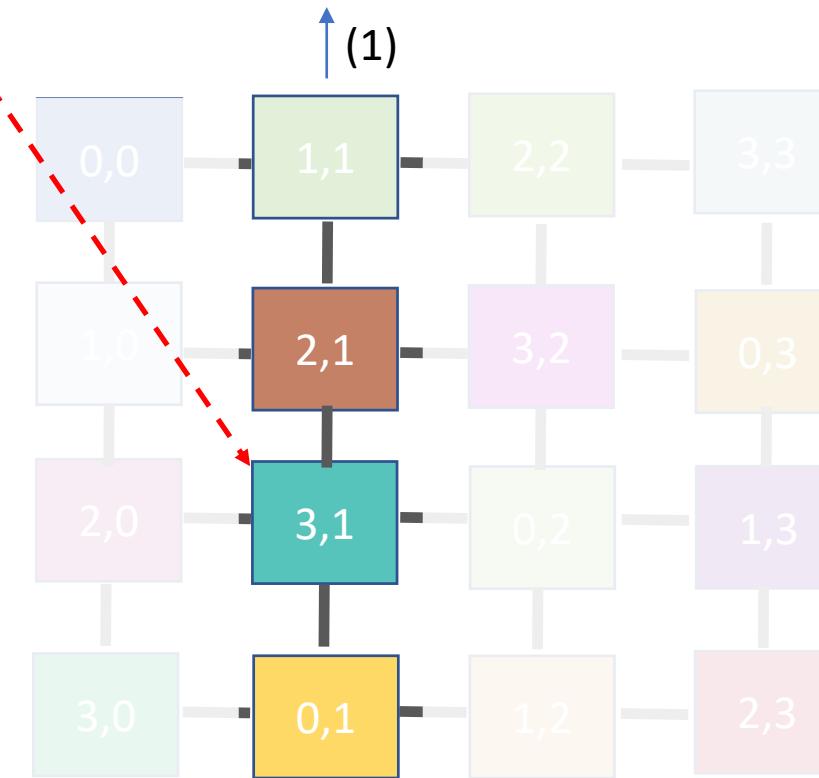
Let us focus on the computation of  $C[2][1]$ .  
Tile  $(2,1)$  will compute it.



A



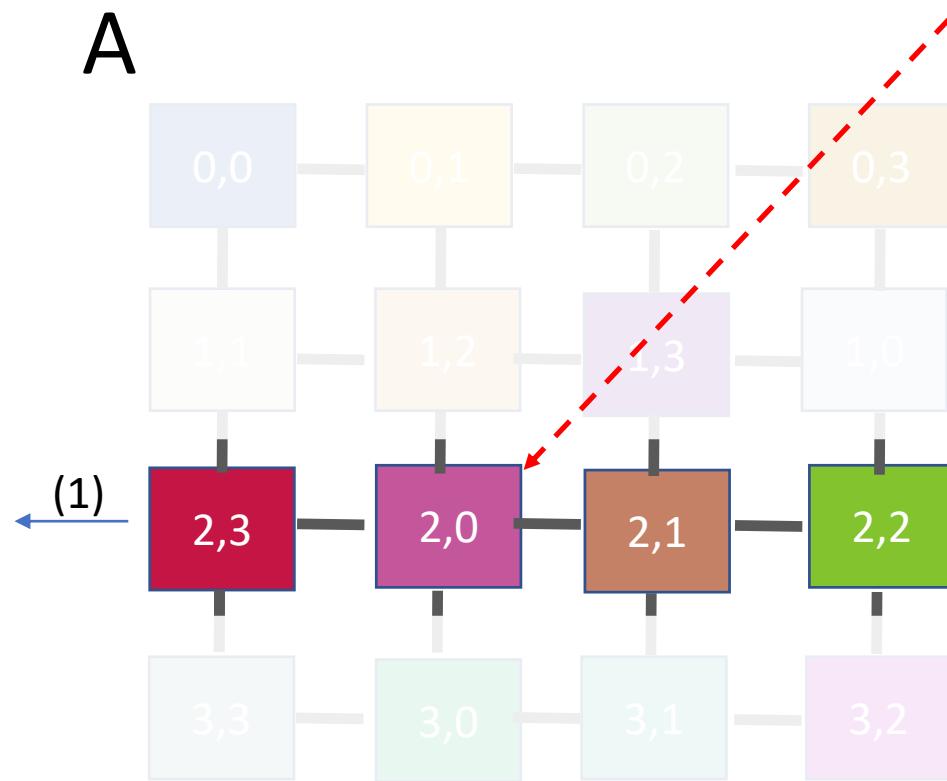
B



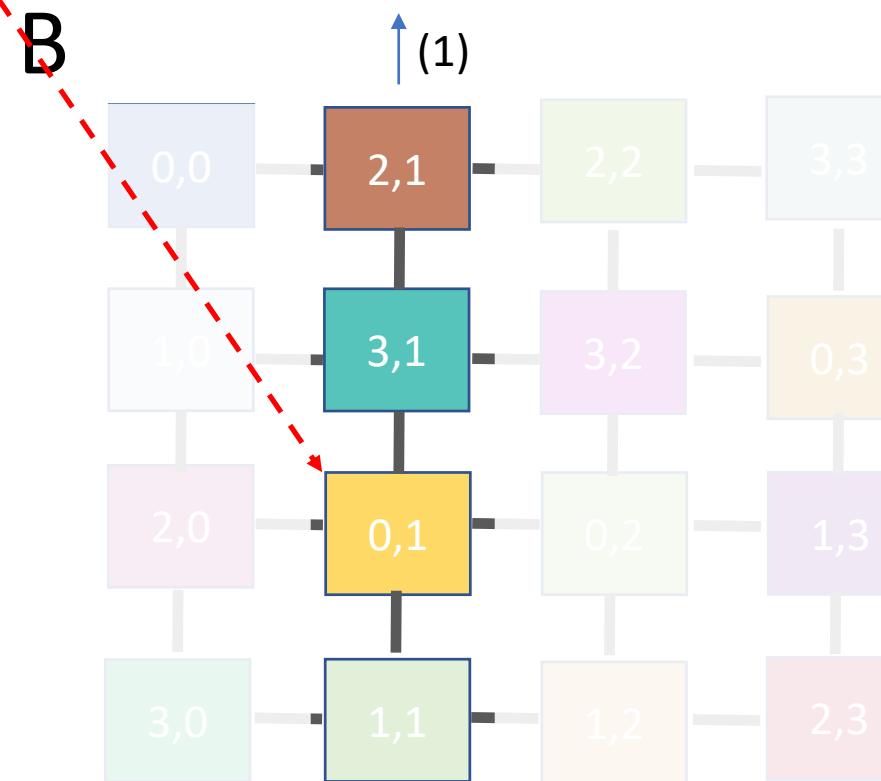
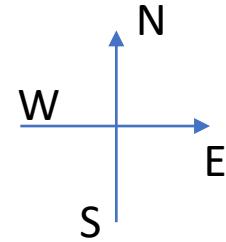
$$C[2][1] = A[2][3] \times B[3][1]$$

STEP 1

# Main Algorithm



Let us focus on the computation of  $C[2][1]$ .  
Tile (2,1) will compute it.

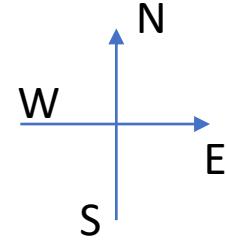


$$C[2][1] = A[2][3] \times B[3][1]$$

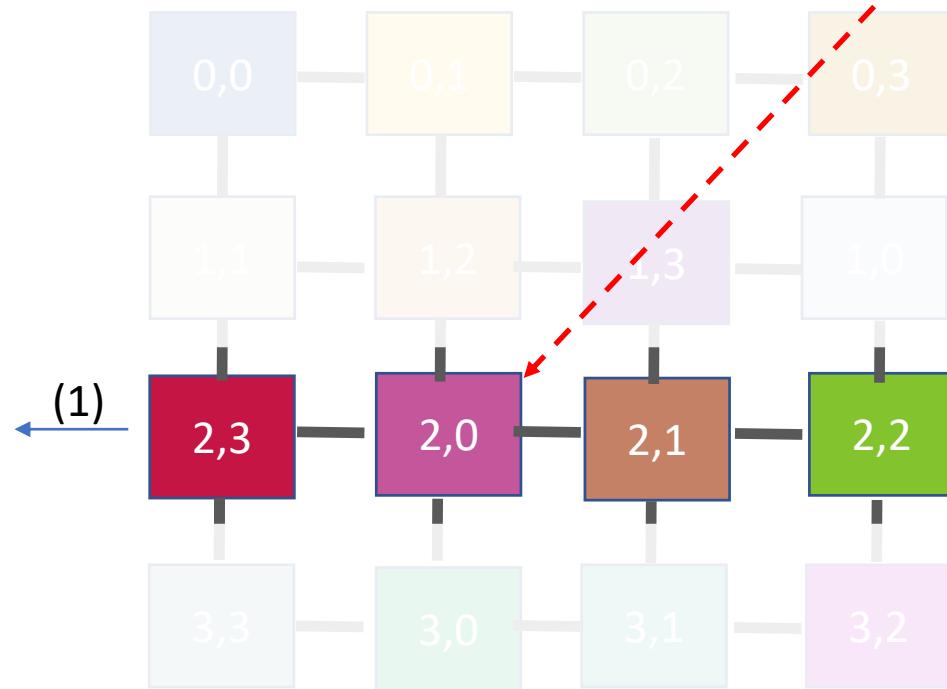
STEP 2

# Main Algorithm

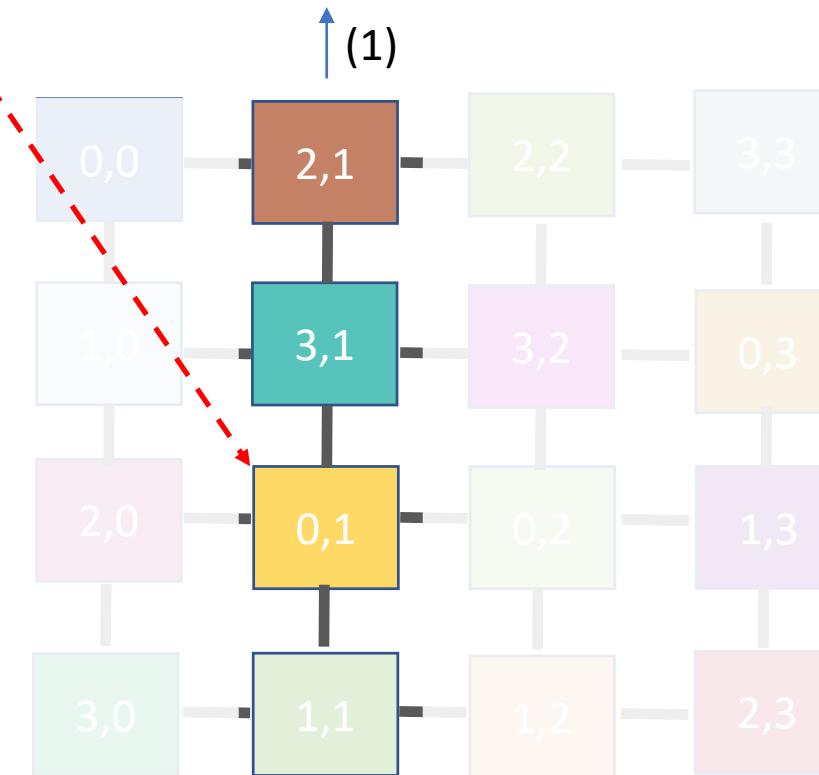
Let us focus on the computation of  $C[2][1]$ .  
Tile (2,1) will compute it.



A



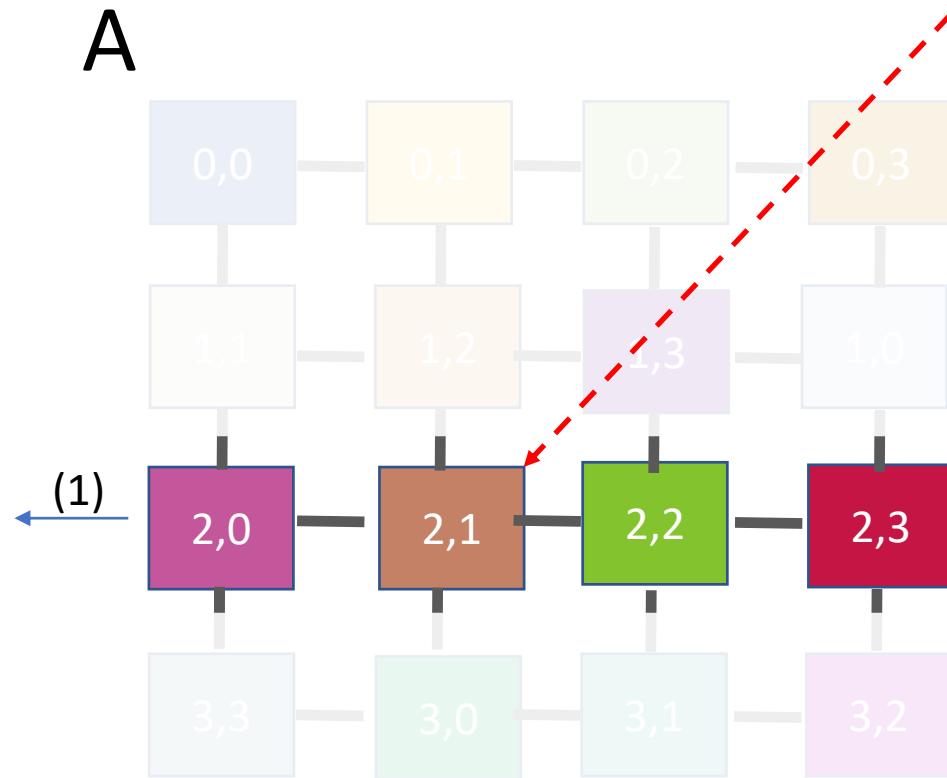
B



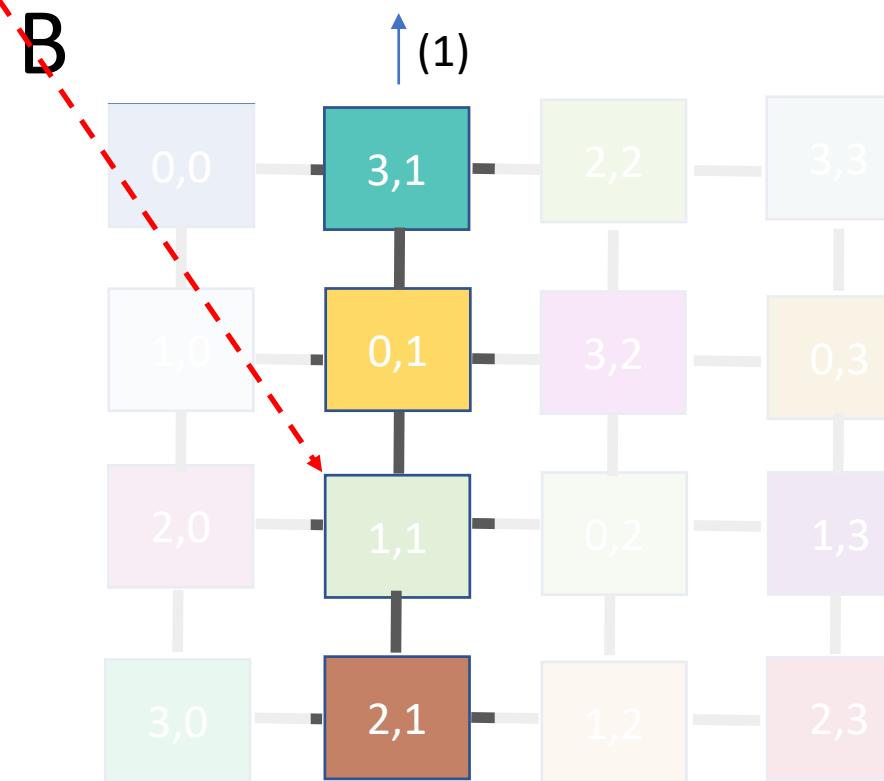
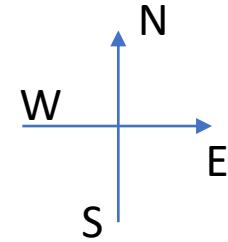
$$C[2][1] = A[2][3] \times B[3][1] + A[2][0] \times B[0][1]$$

STEP 2

# Main Algorithm



Let us focus on the computation of  $C[2][1]$ .  
Tile (2,1) will compute it.

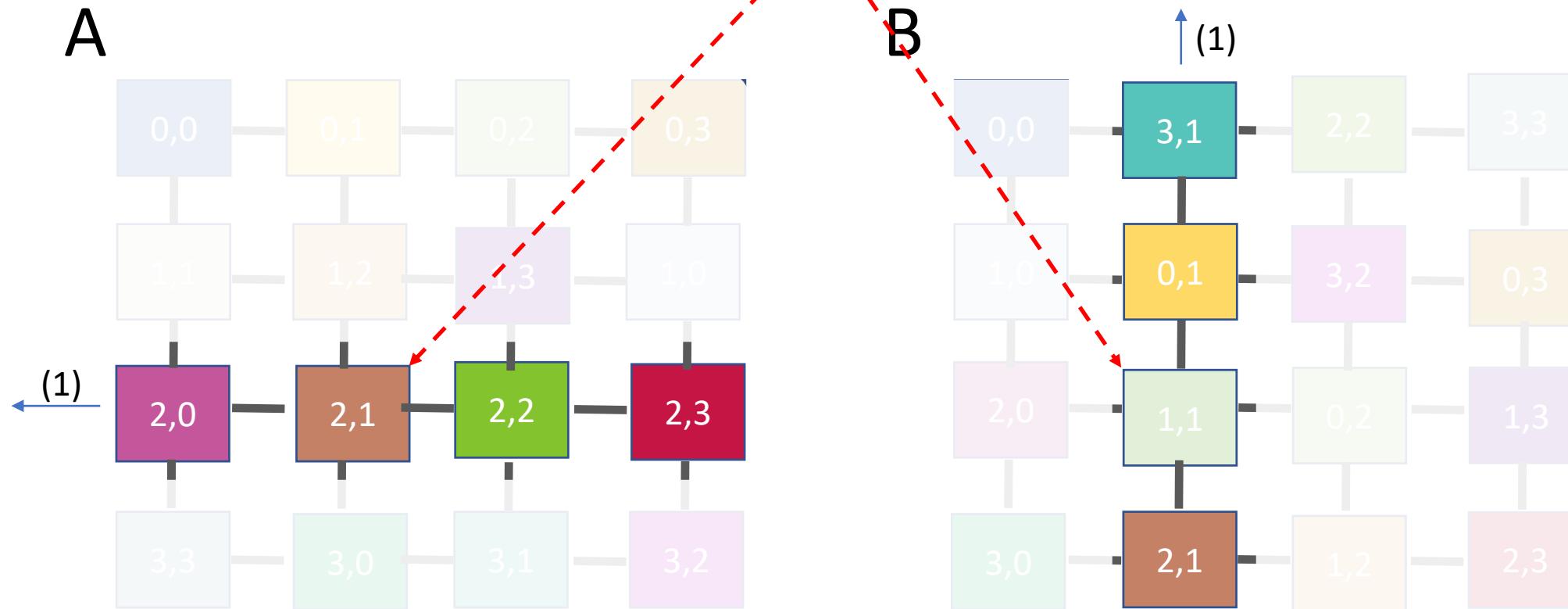
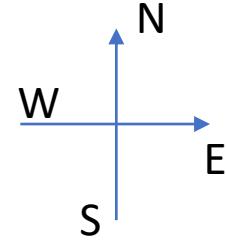


$$C[2][1] = A[2][3] \times B[3][1] + A[2][0] \times B[0][1]$$

STEP 3

# Main Algorithm

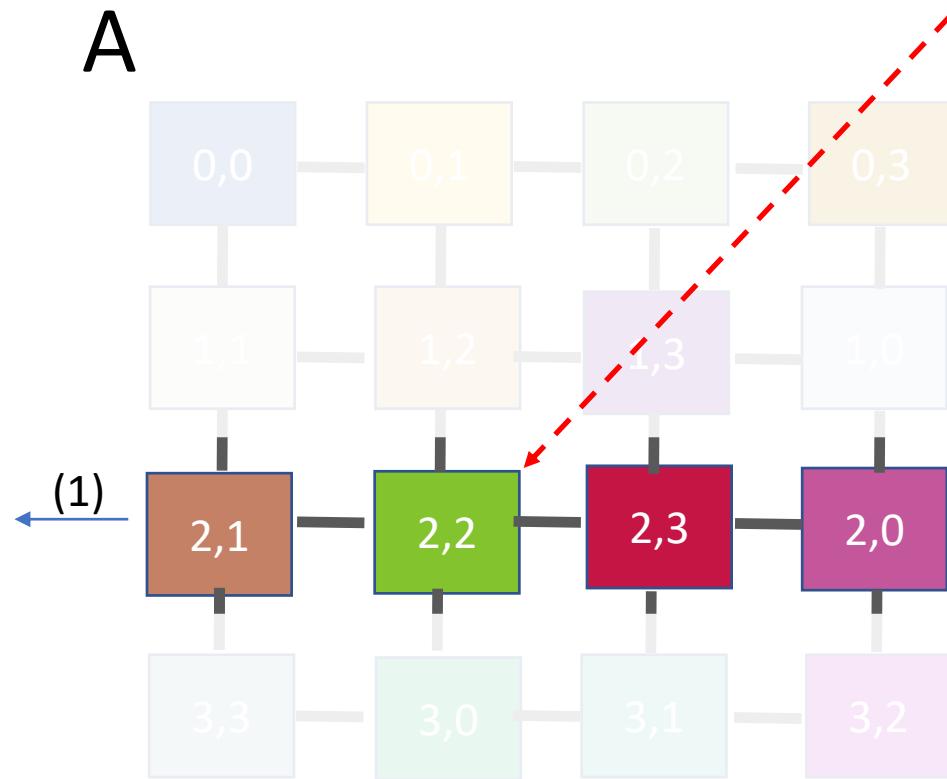
Let us focus on the computation of  $C[2][1]$ .  
Tile  $(2,1)$  will compute it.



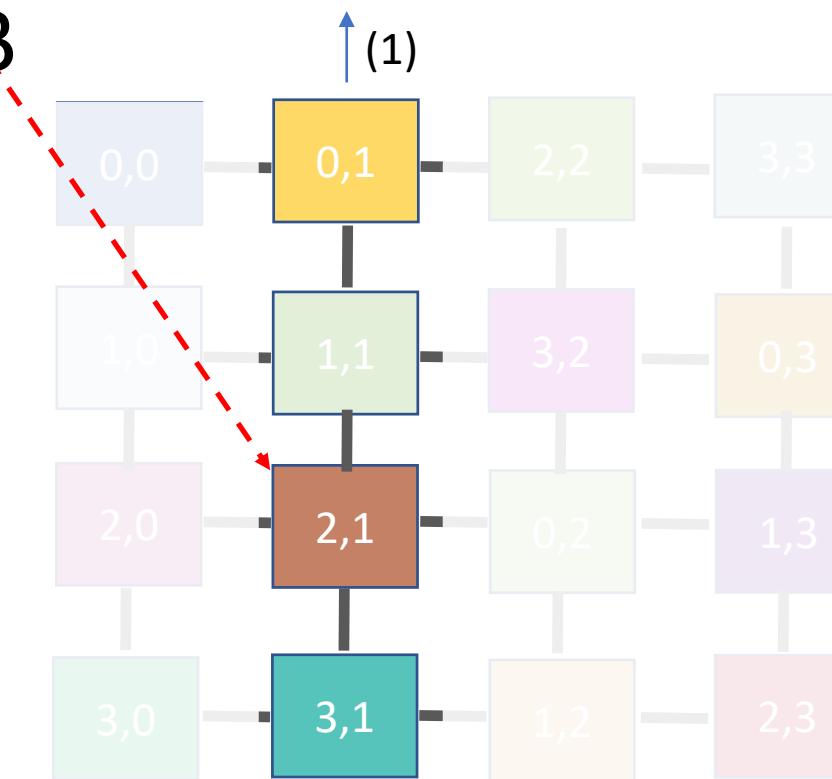
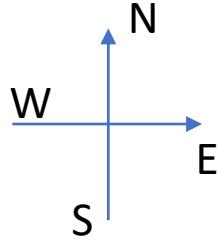
$$C[2][1] = A[2][3] \times B[3][1] + A[2][0] \times B[0][1] + A[2][1] \times B[1][1]$$

STEP 3

# Main Algorithm



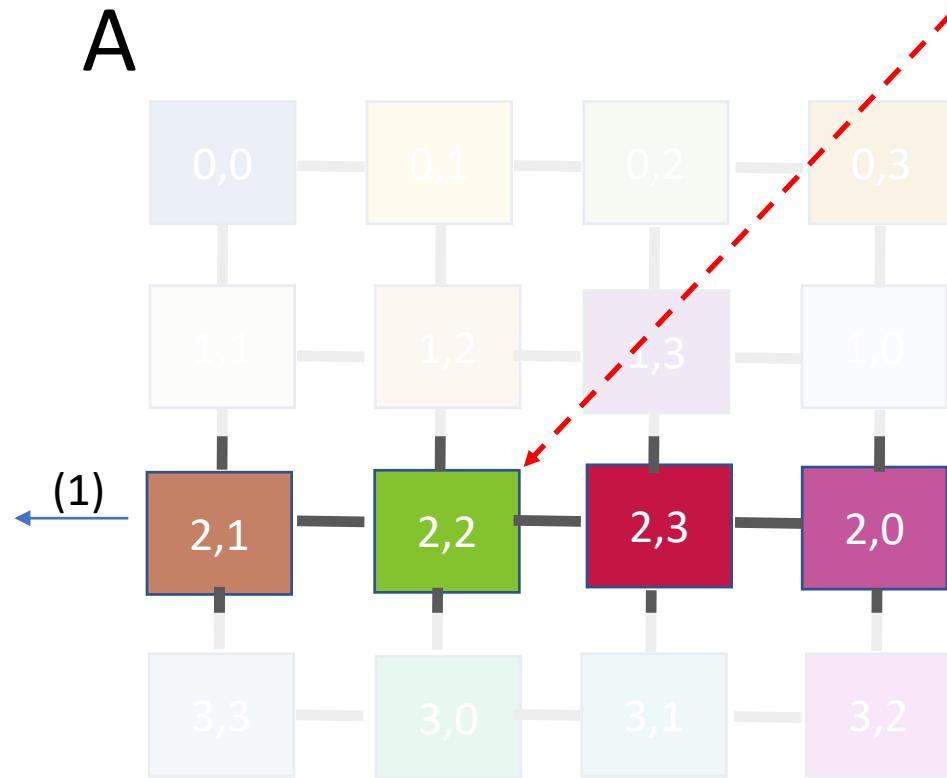
Let us focus on the computation of  $C[2][1]$ .  
Tile (2,1) will compute it.



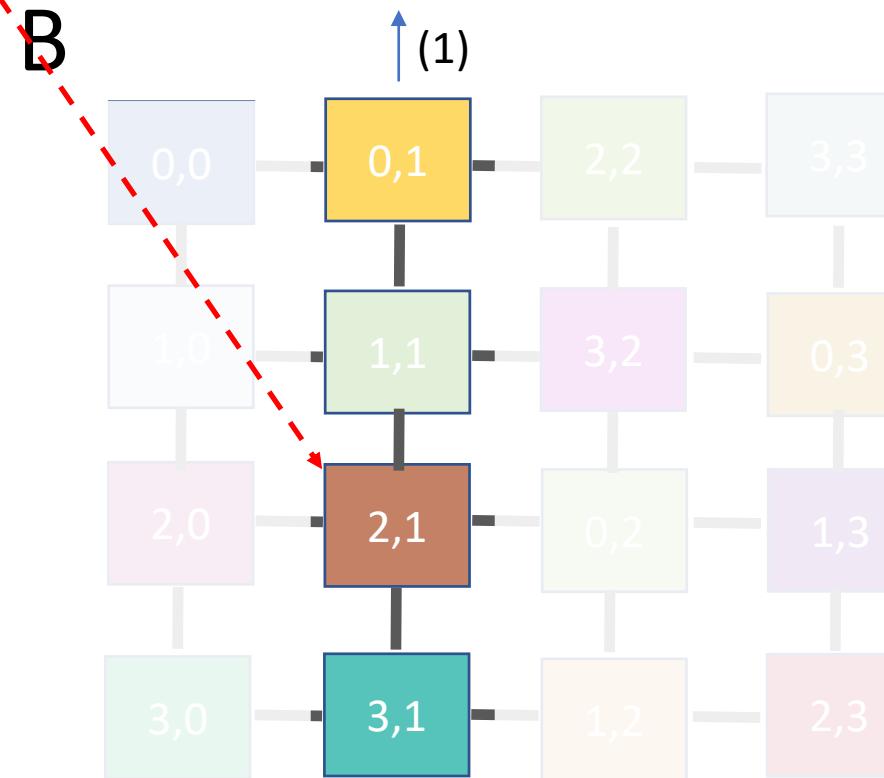
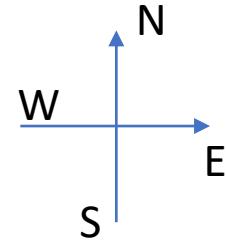
$$C[2][1] = A[2][3] \times B[3][1] + A[2][0] \times B[0][1] + A[2][1] \times B[1][1]$$

STEP 4

# Main Algorithm



Let us focus on the computation of  $C[2][1]$ .  
Tile (2,1) will compute it.



$$C[2][1] = A[2][3] \times B[3][1] + A[2][0] \times B[0][1] + A[2][1] \times B[1][1] + A[2][2] \times B[2][1]$$

Proc (2,1) outputs C[2][1]

STEP 4