

**EE234**

**Microprocessor Systems**

**Final Exam**

**Dec. 15, 2020. (11am – 1:30pm)**

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**Name:**

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| Problem | Points |  |
|---------|--------|--|
| 1       | 20     |  |
| 2       | 20     |  |
| 3       | 30     |  |
| 4       | 20     |  |
| 5-1     | 20     |  |
| 5-2     | 30     |  |
| 6       | 20     |  |
| Total   | 160    |  |

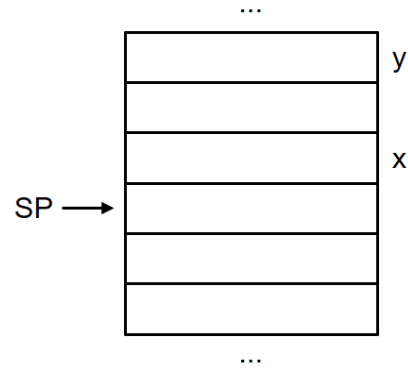
## Problem 1 (1-D Array, 20 points)

All the registers R# are 32-bit registers. “int” is a 32-bit signed integer data type. Write an assembly code for the “for” loop in the following C code. The memory figure shows the stack pointer (SP) and the locations of the variables x and y.

```
int y[10];
int* x = new int[10];

...

for ( int k = 0 ; k < 10 ; k++ )
    x[k] = y[k];
```



```
MOV R0, #0 // k
loop1:
    CMP R0, #10
    BGE finish
    MUL R1, R0, #4 // 4*k
    ADD R2, R1, SP
    LDR R2, [R2, #12] // R2 = y[k]
    LDR R3, [SP, #4] // x
    ADD R3, R3, R1 // &(x[k])
    STR R2, [R3] // x[k] = R2 = y[k]
    ADD R0, R0, #1
    B loop1
finish:
```

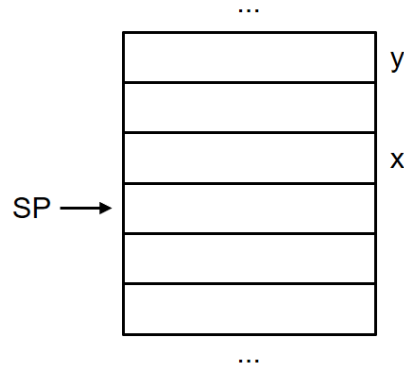
## Problem 2 (1-D Array, 20 points)

All the registers R# are 32-bit registers. “long” is a 64-bit signed integer data type. Write an assembly code for the “for” loop in the following C code. The memory figure shows the stack pointer (SP) and the locations of the variables x and y.

```
long y[10];
long* x = new long[10];

...

for ( int k = 0 ; k < 10 ; k++ )
    x[k] = y[k];
```



```
MOV R0, #0 // k
loop1:
    CMP R0, #10
    BGE finish
    MUL R1, R0, #8 // 8*k
    ADD R2, R1, SP
    LDR R2, [R2, #12] // R2 = LO(y[k]) LO is the lower 4B.
    LDR R3, [R2, #16] // R3 = HI(y[k]) HO is the upper 4B.
    LDR R4, [SP, #4] // x
    ADD R4, R4, R1 // &(LO(x[k]))
    STR R2, [R4] // LO(x[k]) = R2 = LO(y[k])
    STR R3, [R4, #4] // HI(x[k]) = R3 = HI(y[k])
    ADD R0, R0, #1
    B loop1
finish:
```

### Problem 3 (2-D Array, 20 points)

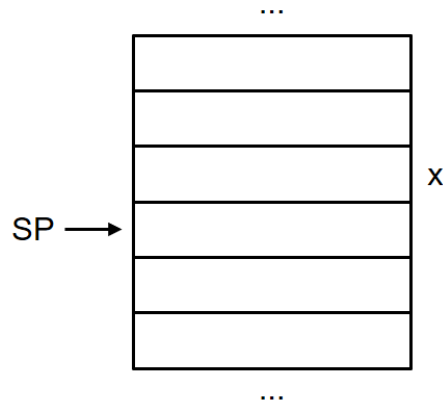
All the registers R# are 32-bit registers. “int” is a 32-bit signed integer data type. Write an assembly code for the nested “for” loop in the following C code. The memory figure shows the stack pointer (SP) and the locations of the variables x and y.

```
int** x = new int*[2];

for ( int k = 0 ; k < 2 ; k++ )
    x[k] = new int[4];

...

for ( int i = 0 ; i < 2 ; i++ ) {
    for ( int k = 0 ; k < 4 ; k++ ) {
        x[i][k] = 0;
    }
}
```



Note: Try to minimize the number of memory access instructions (LDR, STR) executed. However, you should strictly follow the flow of the program (i.e., you should have two nested loops, etc.).

```
MOV R0, #0 // int i = 0
MOV R12, #0 // constant 0
LDR R3, [SP, #4] // x
loop1:
    CMP R0, #2
    BGE finish
    MUL R2, R0, #4 // 4*i
    ADD R4, R2, R3 // &(x[i])
    LDR R4, [R4] // x[i]
    MOV R1, #0 // int k = 0
loop2:
    CMP R1, #4
    BGE loop3
    MUL R2, R1, #4 // 4*k
    ADD R4, R4, R2 // &(x[i][k])
    STR R12, [R4] // x[i][k] = 0
    ADD R1, R1, #1
    B loop2
loop3:
    ADD R0, R0, #1 // i++
    B loop1
finish:
```

## Problem 4 (Estimation of Memory Consumption, 20 points)

Estimate how many bytes are used for the array `x` in the following C code. You should include the memory space used for variable `x` itself. "int" is a 32-bit signed integer data type.

```
int**** x = new int***[3];

for ( int k = 0 ; k < 3 ; k++ )
    x[k] = new int**[4];

for ( int i = 0 ; i < 3 ; i++ ) {
    for ( int k = 0 ; k < 4 ; k++ ) {
        if ( k % 2 == 0 ) {
            x[i][k] = new int*[5];
            for ( int m = 0 ; m < 3 ; m++ )
                x[i][k][m] = new int[6];
        }
    }
}
```

`x` needs 4B.

The first "new" operation reserves 3 int\*\*\* spaces →  $3 \cdot 4B = 12B$ .

The first "for" loop reserves 4 int\*\* spaces for each `x[k]` →  $3 \cdot (4 \cdot 4B) = 48B$ .

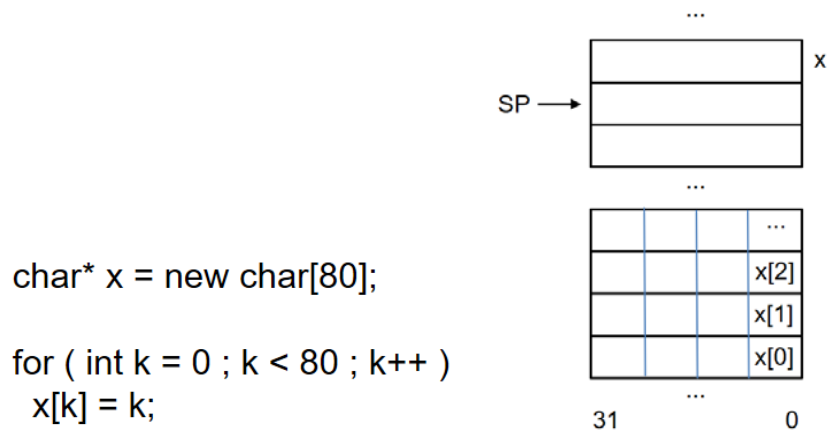
In the nested "for" loop, the "if" statement is executed six times. In the "if" statement, the "new" statement reserves 5 int\* spaces →  $6 \cdot (5 \cdot 4B) = 120B$ .

In the "for" loop in the "if" statement, the "new" operation reserves 6 int spaces. →  $6 \cdot (3 \cdot (6 \cdot 4B)) = 432B$ .

Thus, total 616 Bytes.

## Problem 5-1 (Array Manipulation I, 20 points)

The “char” data type in C is used to represent 1 byte. If you need an array of  $M$  char-type variables, you will ideally need  $M$  bytes. However, all the memory addresses for LDR and STR instructions should be integer multiples of 4 in the 32-bit ARM architecture (so, for example, you cannot use 0x0001 for a target memory address). Now, let’s take a look at the following C code. It reserves memory space for 80 characters, so ideally it should reserve 80 Bytes in the heap memory. However, it requires some bit manipulations. Thus, a compiler can reserve 320 Bytes in the heap memory and use only the least significant 1B in each word for each  $x[k]$  as follows.



Write an assembly code for the “for” loop in the C code shown above. The memory management should be the same as the compiler above.

```

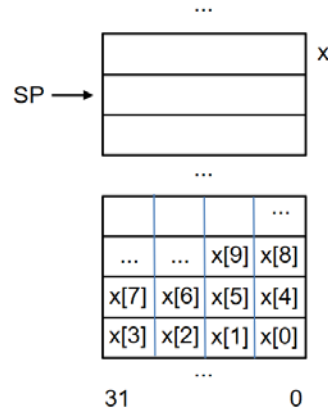
LDR R0, [SP, #4]    // R0 = x
MOV R1, #0         // k = 0
loop1:
CMP R1, #80
BGE finish
MUL R2, R1, #4     // 4*k
ADD R2, R0, R2     // &(x[k])
STR R1, [R2]       // x[k] = k
ADD R1, R1, #1
B loop1
finish:
        
```

## Problem 5-2 (Array Manipulation II, 30 points)

For the C code in Problem 5-1, a different compiler reserves exactly 80 Bytes in the heap space as follows.

```
char* x = new char[80];
```

```
for ( int k = 0 ; k < 80 ; k++ )
    x[k] = k;
```



Write an assembly code for the “for” loop in the C code shown above. The memory management should be the same as the new compiler explained above.

```
LDR R0, [SP, #4]    // R0 = x
MOV R1, #0          // k = 0
loop1:
  CMP R1, #80
  BGE finish
  MOV R2, R1, LSR #2 // k/4
  MUL R2, R2, #4     // 4*(k/4)
  ADD R2, R0, R2
  LDR R3, [R2]       // load x[a+3:a]
  AND R4, R1, #0x3   // k%4
  MUL R4, R4, #8     // # bits to shift to the left
  MOV R5, R1, LSL R4 // shift k to the left
  MOV R6, #0xFF
  MOV R6, R6, LSL R4 // shift 000...011111111 to the left
  LDR R7, =#0xFFFFFFFF
  EOR R6, R6, R7     // inversion of R6
  AND R6, R3, R6     // finally we get x[a+3] x[a+2] 00000000 x[a] (assuming x[a+1]=k)
  ORR R6, R6, R5     // x[a+3] x[a+2] k x[a]
  STR R6, [R2]
  ADD R1, R1, #1     // k++
  B loop1
finish:
```

## Problem 6 (C, 20 points)

All the registers R# are 32-bit registers. “int” is a 32-bit signed integer data type and “long” is a 64-bit signed integer data type. The following table shows the main memory.

```
int** x = new int*[a];

for ( int i = 0 ; i < a ; i++ )
    x[i] = new int[b];
```

“a” and “b” are some constants. Currently, the value of x is 0x4000 as shown in the figure.

- What is the value of \*((int\*) x)? **0x400C**
- What is the value of \*((long\*) x)? **0x4014 400C**
- What is the value of x[2]? **0x4024**
- What is the value of x + 3? **0x400C**
- What is the value of (x[0]+2)? **0x4014**
- What is the value of x[1][2]? **0x4020**
- int\* y = x[1]. What is the value of y[3]? **0x4024**
- long\*\* y = (long\*\*) x. What is the value of y[1]? **0x4014**
- What is the value of &(x[2])? **0x4008**
- long\* y = (long\*) x[0]. What is the value of y[1]? **0x401C 4018**

| Address | Data      |
|---------|-----------|
|         | 31      0 |
| 0x8000  | 0x4000    |
| ...     |           |
| 0x402C  | 0x4000    |
| 0x4028  | 0x402C    |
| 0x4024  | 0x4028    |
| 0x4020  | 0x4024    |
| 0x401C  | 0x4020    |
| 0x4018  | 0x401C    |
| 0x4014  | 0x4018    |
| 0x4010  | 0x4014    |
| 0x400C  | 0x4010    |
| 0x4008  | 0x4024    |
| 0x4004  | 0x4014    |
| 0x4000  | 0x400C    |