Computer Systems

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Primary Hardware Components

- CPU (Central Processing Unit)
- Main memory
- Disk drives
 - Hard disk drive (HDD)
 - Solid state drive (SSD)
- o Input
 - Keyboard, mouse, ...
- Output
 - Monitor, printer, ...





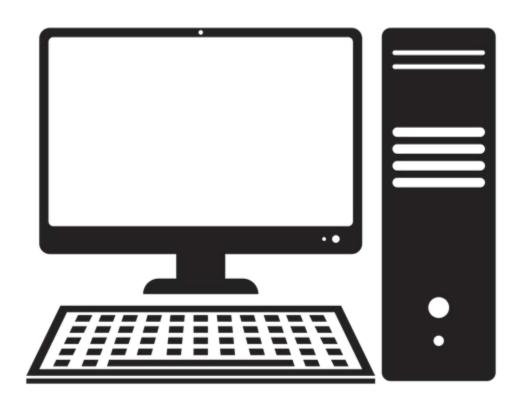




Primary Software Components

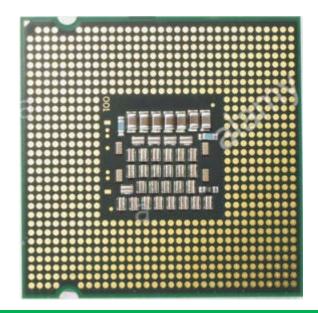
- o Operating System
- o Compiler
- o Linker
- o Loader
- Application

How Each Component Works

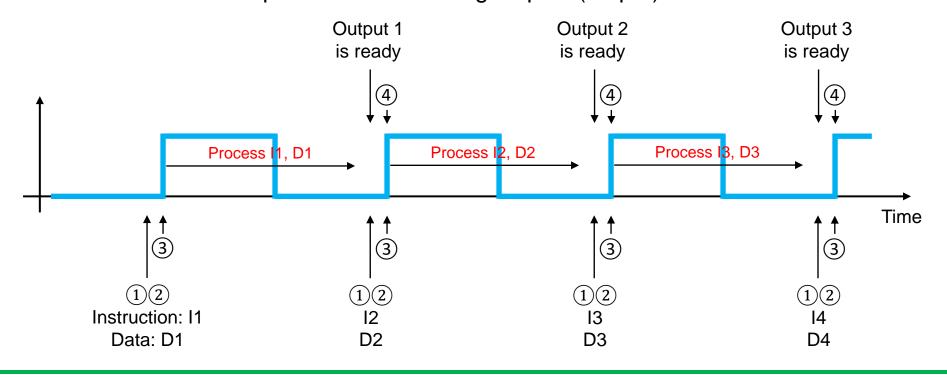


o Pins

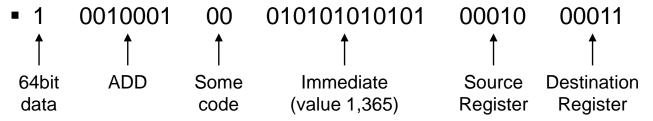
- Input (signals going into the CPU)
- Output (signals coming out from the CPU)
- Power (VDD, GND)
 - VDD: High voltage (e.g., 1V)
 - GND: Low voltage (e.g., 0V)



- o A very simple, conceptual model for CPU operations
 - 1. Place an input data on its data signal pins (input).
 - 2. Place an instruction on its instruction signal pins (input).
 - 3. Provide a clock (then the CPU will run the instruction).
 - 4. Read the data placed on its data signal pins (output).



- O How do the instructions and data look like?
 - 10010001000101010101010001000011
- O What is it?

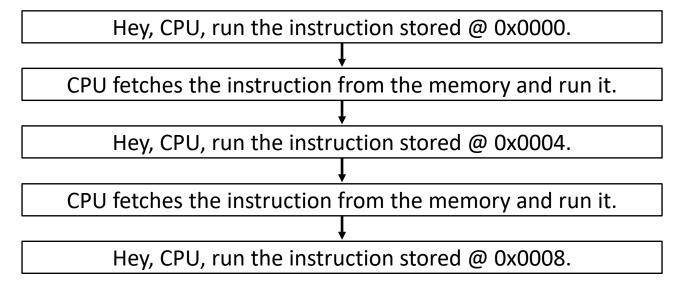


- Add 1,365 and the value stored in the source register.
 - (A register is a temporary memory space. We will study that later.)
- Store the result in the destination register.
- o If an instruction is given, how can we find out what it means?
 - 000001001101011111110001110000011

- Instruction Set Architecture (ISA)
 - The ISA of a CPU defines
 - All the instructions the CPU can perform (addition, subtraction, ...)
 - The formats of the instructions and their machine codes (0110001...)

```
ADD (immediate)
    Add (immediate) adds a register value and an optionally-shifted immediate value, and writes the result to the destination register.
    This instruction is used by the alias MOV (to/from SP).
31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
sf 0 0 1 0 0 0 1
                           != 1x
                                                  imm12
                                                                                                 Rd
   op S
                            shift
32-bit (sf == 0)
    ADD <Wd|WSP>, <Wn|WSP>, #<imm>{, <shift>}
64-bit (sf == 1)
    ADD <Xd|SP>, <Xn|SP>, #<imm>{, <shift>}
    integer d = UInt(Rd);
    integer n = UInt(Rn);
    integer datasize = if sf == '1' then 64 else 32;
    bits(datasize) imm;
    case shift of
        when '00' imm = ZeroExtend(imm12, datasize);
        when '01' imm = ZeroExtend(imm12:Zeros(12), datasize);
        when '10' SEE "ADDG, SUBG";
        when '11' ReservedValue();
```

- Yes, we can run a program like that.
- o However, we want to automate the whole process. How?
 - Instructions, data: Store them in memory.
 - Tell the CPU to fetch the instruction stored at a certain memory address and run it.
 - Then, tell the CPU where the next instruction is.



Hardware – Memory (DRAM)

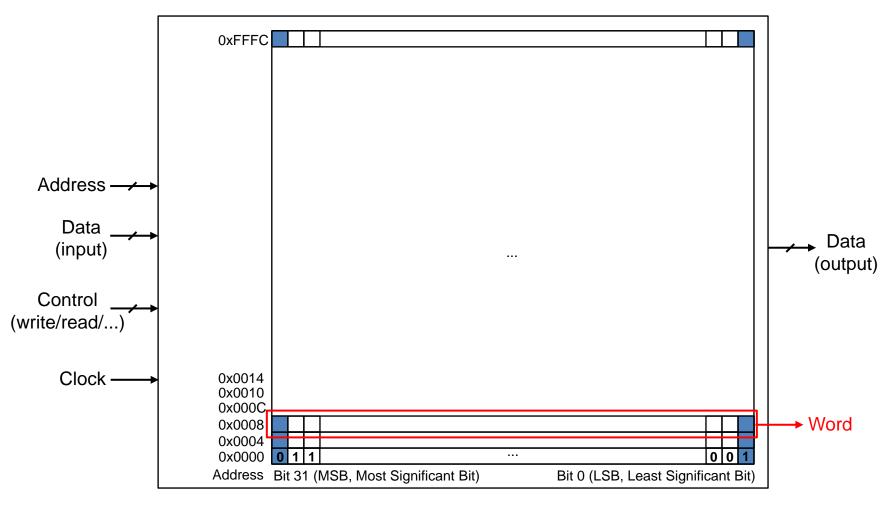
o Pins

- Input (signals going into the memory)
- Output (signals coming out from the memory)
- Power (VDD, GND)
 - VDD: High voltage (e.g., 1V)
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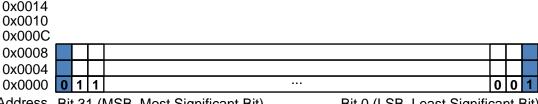
Hardware – Memory

o Abstract model



Hardware – Memory

- How to calculate the size of a given memory chip
 - Word (data) width: w (bits)
 - # words: k
 - Size: $k \cdot w$ (bits) = $\frac{k \cdot w}{g}$ (Bytes)
 - Example
 - w = 32
 - $k = \frac{2^{16}}{4} = 2^{14}$
 - $s = 2^{19}$ bits = 2^{16} Bytes = 2^6 KBytes = 32 KBytes
- If the address bus has 32 bits, what's the max. memory capacity it can support?
 - Answer: 4GB

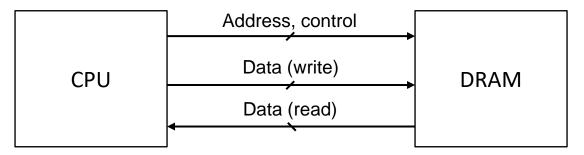


Address Bit 31 (MSB, Most Significant Bit)

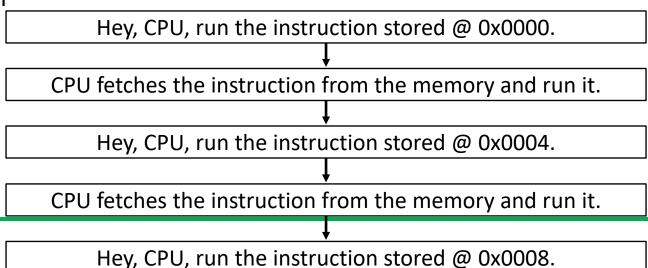
Bit 0 (LSB, Least Significant Bit)

Hardware – CPU and Memory (DRAM)

- Now, we know how a CPU chip and a memory chip look like.
- Let's connect them.



- o Finally, we can automatically run a program.
 - Load your program and data into the main memory.
 - Then



Hardware – CPU and Memory (DRAM)

- CPUs are fast (high-speed).
 - Running an instruction takes one clock cycle (~300ps).
- Memory is slow (low-speed).
 - Loading a data takes several hundreds of clock cycles.
- Question
 - Why can't we integrate a large amount of memory in a CPU?
- Solution
 - Let's have a small, but fast memory inside a CPU.

- o Register File (RF)
 - A register is a memory space that can store a word.
 - A register file is a set of registers.
- The input/output behavior of an RF is very similar to that of the abstract model of a general memory chip.
- Since RF should be very fast, RF has only a few registers (16~32).

R15
R14
R13

R1	
R0	

- Comparison with vs. without RF
 - Instructions: Add data1 (stored @ memory address 0x4000), data2 (@ 0x4004), data3 (@ 0x4008), and data4 (@0x400C), then store the result @ 0x5000.
 - Let's assume that a memory instruction takes 200 clock cycles.

Load data @ 0x4000	+200 cycles
Load data @ 0x4004	+200
Add the two	+1
Store @ 0x8000	+200
Load data @ 0x4008	+200
Load data @ 0x400C	+200
Add the two	+1
Store @ 0x8004	+200
Load data @ 0x8000	+200
Load data @ 0x8004	+200
Add the two	+1
Store @ 0x5000	+200

Load data @ 0x4000 to R1	+200
Load data @ 0x4004 to R2	+200
Add the two and store to R3	+1
Load data @ 0x4008 to R1	+200
Load data @ 0x400C to R2	+200
Add the two and store to R4	+1
Add R3, R4 and store to R1	+1
Store R5 @ 0x5000	+200

Hardware – CPU and Memory (DRAM)

Summary

- A CPU has a Register File (RF) consisting of 16~32 registers.
- To perform a task in a CPU, you should
 - Load the data from the memory
 - Store it in the RF
 - Perform a task (addition, subtraction, ...) on the registers
 - Store the result back to the RF
 - Store the result back to the memory (if necessary)
- RFs are small, RFs should be used only for temporary space.
- All the instructions and data should be stored in the main memory.
 - We will study this later.
 - For now, we will assume that a program and its data are stored in the memory.
- A CPU has an instruction set architecture (ISA).
 - We will study this later.
- How to run a program automatically
 - Hey, CPU, the next instruction is @ X. Do it. Then, the next instruction is @ Y. ...

Software – Application Programming

- Let's make a small program that adds two values stored @ 0x4000 and @ 0x4004 and store the result @ 0x5000.
- o How would you do it?
 - Make a pseudo code.

Load data @ 0x4000 to R1 Load data @ 0x4004 to R2 Add the two and store to R3 Store R3 @ 0x5000

Translate it into a machine code. (You need an ISA reference for this.)

Load data @ 0x4000 to R1 Load data @ 0x4004 to R2 Add the two and store to R3 Store R3 @ 0x5000



Now, it is very error-prone, inefficient, and hard to maintain.

Software – Assembly

- An assembly language is a human-friendly programming language that makes the machine code programming much easier.
- An assembly code has a kind of one-to-one correspondence with a machine code.

Assembly Machine code

LDR R1, #0x4000
(Load data @ 0x4000 to R1) 00011000100110110001011111000

- The translator (from an assembly code to a machine code) is an assembler.
 - The backward translator is a disassembler.

Software – Application Programming

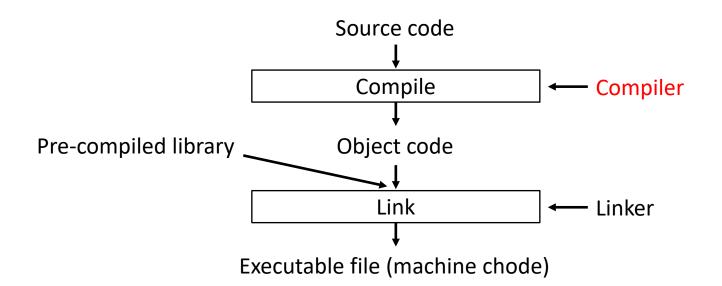
- Yes, coding with an assembly language looks much easier than coding directly with the machine language.
- However, should we need to remember all the instructions and the syntax?
 - Yes, we do...
 - But we can remember only 20~30 instructions.
 - ADD (addition)
 - SUB (subtraction)
 - MUL (multiplication)
 - AND (logical AND)
 - OR (logical OR)
 - EOR (logical exclusive-OR)
 - LSR/LSL (logical shift right/left)
 - MOV (move data)
 - CMP (compare)
 - LDR/STR (load/store)

Software – High-Level Programming Languages and Compilers

- Unfortunately, assembly programming is still extremely hard, timeconsuming, error-prone, unportable, and human-unfriendly.
- High-level programming languages (C/C++, ...)
 - Easy
 - Efficient
 - Human-friendly
 - High-performance
 - Portable (independent of CPUs)
 - Different CPUs have different ISAs.
 - Assembly languages are dependent on CPUs.

Software – High-Level Programming Languages and Compilers

Application programming

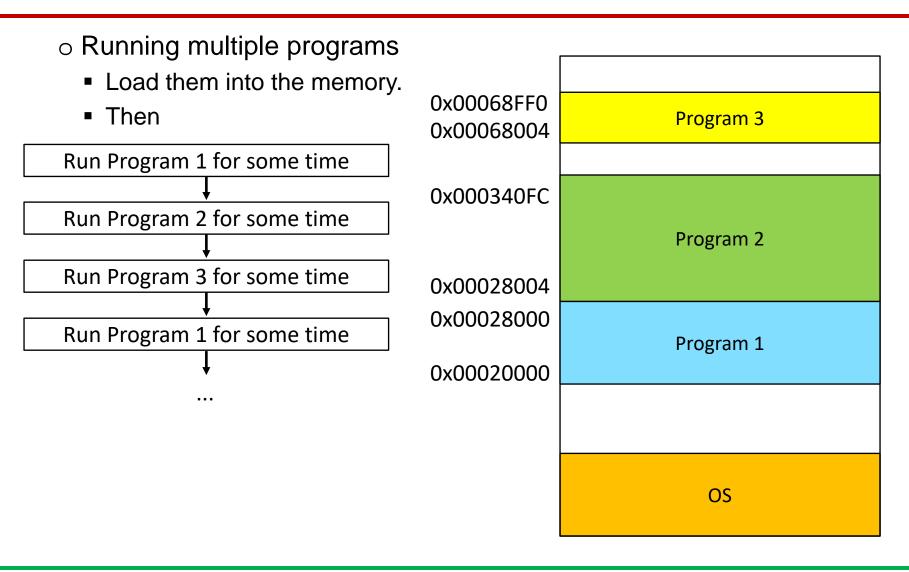


o If you run a single program, managing the program is easy. Just load your program to the main memory starting from 0x0000 and run it from 0x0000.

0x01C8	В
0x01C4	BNE
	•••
8000x0	ADD
0x0004	LDR
0x0000	LDR

- What if you want to
 - run multiple programs at the same time?
 - manage input/output devices
 - manage the main memory dynamically (on-the-fly)
 - **-** ...
- Operating system (OS)
 - Process management
 - Memory management
 - Storage management
 - I/O device management
 - File system
 - Security
 - ..

- An I/O device has a special communication protocol.
 - OS has device drivers.
 - Application programmers don't need to develop device drivers.
- Multiple programs ask for resources (storage, memory, CPU, ...).
 - OS manages them intelligently to maximize the system performance, usability, etc.
- o How does an OS manage the main memory?



- An issue when you run multiple programs
 - Some codes use relative memory addresses.
 - If Instr 1 is located at 0x0000
 - Instr 6 is @ 0x0014
 - Instr 9 is @ 0x0020
 - If Instr 1 is located at 0x4000
 - Instr 6 is @ 0x4014
 - Instr 9 is @ 0x4020

Instr 10	SUB		
Instr 9	ADD		
Instr 8	ADD		
Instr 7	SUB		
Instr 6	ADD	←	
Instr 5	Go to Instr 20		
Instr 4	Go to Instr 9 if]	
Instr 3	Go to Instr 6 if		
Instr 2	LDR		
Instr 1	LDR		

The relative memory addresses are determined when the program is loaded into the memory.

Software - Loader

Solution

- The compiled machine code has relative memory addresses.
- When the program is loaded into the memory, fill in the relative addresses.
 - Who does it? Loader

Instr 10	SUB	
Instr 9	ADD	
Instr 8	ADD	
Instr 7	SUB	
Instr 6	ADD	
Instr 5	Go to Instr 20	
Instr 4	Go to Instr 9 if	
Instr 3	Go to Instr 6 if	
Instr 2	LDR	
Instr 1	LDR	

0x4064	SUB	
0x4060	ADD	
0x405C	ADD	
0x4058	SUB	
0x4054	ADD	
0x4050	Go to 0x408C	
0x404C	Go to 0x4060 if	
0x4048	Go to 0x4054 if	
0x4044	LDR	
0x4040	LDR	

Software

- o Summary
 - Machine language: 00011000001010101101100001010101
 - Assembly language: ADD R1, R2, R3
 - High-level programming language: a = b + c;
 - Compiler
 - Linker
 - Loader
 - Operating system