

Chapter 4

- 50+ years of CS and SE devoted to issues of composability: how to create solutions to big problems by combining (composing) solutions to smaller problems
 - Mechanisms e.g. intrinsic locks; GC
 - Techniques Invariants, pre- and postconditions; confinement; delegation; etc. (Ch4)
 - Libraries working code embodying the techniques for specific domains (Ch 5)
- This chapter mainly about techniques



Invariants – <u>the</u> fundamental technique

- What property is always true of an object when it is in a correct state?
- For sequential programming, the class invariant gives you critical information about what each public method needs to achieve
- For concurrent programming, the class invariant tells you which fields are related and therefore have to be protected by a single lock



Post-conditions and Pre-conditions

- The post-condition of a method tells you what that method is supposed to accomplish
 - (A no-op will preserve the invariant but that's not very interesting or useful!)
- The pre-condition tells you what (beyond the invariant) is supposed to be true for a method to successfully reach the post-condition
 - In sequential programming calling a method when its precondition is false is an error
 - In concurrent programming we can wait for the precondition to become true



Two common problems

- Composing a thread-safe class from unsafe building blocks
- Composing a thread-safe class when the building blocks are already thread-safe



Confinement Technique (4.2) – threadsafe object built from unsafe objects

 Allow access to a thread-unsafe object only through another object that is thread-safe

```
public class PersonSet {
    private final Set<Person> mySet = new
    HashSet<Person>();
    public synchronized void add(Person p) {
        mySet.add(p); }
    public synchronized boolean contains(Person p) {
        return myset.contains(p); }
```

- }
- hashSet is not ThreadSafe, PersonSet is
- Idea of ownership: PersonSet owns mySet but probably not the Persons contained in it



Danger in Confinement Technique

 Inadvertant publication of what is supposed to be private (confined) mutable state

public synchronized MutablePoint getLocation(String id) {
 MutablePoint loc = locations.get(id);

```
return loc == null ? Null : new MutablePoint(loc);
```

```
}
```

```
Public synchronized setLocation(String id, int x, int y) {
   MutablePoint loc = locations.get(id);
   if (loc == null) { ... exception ...}
   loc.x = x; loc.y = y
```

}

• My preference would be to express this interface using ImmutablePoints.



Thread-safe objects built from thread-safe components – Delegating safety (4.3)

- Delegation: giving responsibility for thread safety to the object(s) containing this object's state
 - ConcurrentMap (TS) instead of Map (not TS)
 - Atomic<foo>
- If this object's state involves multiple other objects delegation may or may not work
 - If the sub-objects are independent, ok
 - If the sub-objects are related, this object must provide its own synchronization – even if all the sub-objects are themselves thread-safe



Example

```
class PongPaddle {
   private final AtomicInteger left = new AtomicInteger(0);
   private final AtomicInteger right = new AtomicInteger(1);
   public void move(int dx) { left.getAndAdd(dx);
   right.getAndAdd(dx); }
   public void changeWidth(int dw) { right.getAndAdd(dw); }
   public boolean hit(int pos) {
      return left.get()<=pos && pos <= right.get();
    }
}</pre>
```

- No visibility concerns
- What is the invariant that relates left and right?
- What should we do to fix it?



Reduce, Reuse, Recycle

- Don't Repeat Yourself (DRY)
- I am very anti cut-and-paste coding
 - Hard on the reader
 - Hard on the maintainer
 - Instead of 1 change, n changes
 - Instead of 1 bug, n bugs
- Design code so there only needs to be one copy (use parameterization, polymorphic parameterization)
- Even better, reuse existing code that does almost the right thing
- How does this interact with synchronization?



Adding functionality (4.4)

- Example: add putlfAbsent to a collection that already supports atomic contains() and add() methods
- Four approaches
 - Modify existing class
 - Extend existing class or
 - Wrap existing class "client-side" locking
 - Composition -



1. Modify existing class

- Assuming the existing class is already thread-safe:
 - Introduce a new method that uses the same synchronization technique already in use
 - Best way but
 - Assumes you have control over the existing class
 - Would not be the case for library classes



2. Extend the existing class

```
public class BetterVector<E> extends Vector<E> {
    public synchronized boolean putIfAbsent(E x) {
        boolean absent = !contains(x);
        if (absent) { add(x); }
        return absent;
    }
}
```

- Vector is a thread-safe library class
- Vector provides enough primitive building blocks to allow construction of putIfAbsent
- Fig. 4.13 Note the benefit of re-entrant locks!
- Note the implicit assumption that we understand the way that Vector does synchronization – using intrinsic locks, in this case
- Note that we don't have any dependency on Vector's implementation



3. Client-side locking

- Assume v is a thread-safe list obtained from
 v = Collections.synchronizedList(new ArrayList<E>());
- Type of this object is List<E> -- not extendable
- Any code that wants to do putIfAbsent item x to such a list, v, can write synchronized (v) {

```
if (!v.contains(x)) v.add(x);
```

- }
- Could be placed in a helper class beware you have to synchronize on the *list* and not on the helper object – see Figs 4.14 and 4.15
- Still depending on knowing the synch policy for the wrapped object
- ... and spreading the knowledge about the synchronization policy far and wide



Composition

- Mimic the idea of Collections.synchronizedList
 - Provide all the synchronization in a new object that extends the functionality of an existing object instance (not class)
 - Delegates most operations to the existing object

```
Public class ImprovedList<T> implements List<T> {
    private final List<T> list;
    public ImprovedList(List<T> list) {
        this.list = list; }
    public synchronized boolean putIfAbsent(T x) ...
    public synchronized boolean contains(T x) {
        return list.contains(x); }
```

•••

Note this is similar to how we handled an non-thread-safe object



Intro to Chapter 5 – Building Blocks

- Chapter 4 was about low-level techniques
- This chapter is about libraries embodiments of the techniques
- Section 5.1 Synchronized collections read to see why you want to use Concurrent collections instead
 - The idioms described are even more unsafe than asserted in the book because of visibility problems



Chapter 5 topics

- Concurrent collections (5.2)
- The ubiquitous producer-consumer pattern (5.3)
- Interruptable methods (5.4)
- Primitive synchronizers (5.5)