

Questions about Assignment

- Interpreting error messages:
- Error in process <0.38.0> with exit value: {function_clause,[{ring,waitSend,[<0.39.0>]}]}
- Exit values are documented in the Erlang manual online.
- This one means: died trying to call ring:waitSend with argument <0.39.0> (a Pid) because there was no matching function clause. Problem: the argument name began with a lowercase letter.



Objects simulated in C

- In C, objects can be emulated by using structs
 typedef obj = struct {mltype *m1; m2type *m2;
 dltype d1... }
 obj *o1 = oFactory(...init values);
 o1->ml(o1, ...);
- C function pointers refer only to function code; not function closures (code together with lookup environment) so all data needed to specialize the code has to be included in the struct
- Conceptually not hard, but requires explicit code to do it



Objects simulated in functional style

- Objects can be simulated using records that contain function closures
- Example: single method that returns previous state and sets new state:

```
newFoo(InitState) -> (fun(NewState)->
```

```
{InitState, newFoo(NewState)}
end).
```

```
F0 = newFoo(3).
```

```
\{V1, F1\} = F0(hi).
```

```
\{V2, F2\} = F1(6).
```



Generalizing

• Now both result and new state are functions of previous state and args.

```
newFoo(InitState) -> (fun(Args)->
```

```
{f1(InitState, Args),
newFoo(f2(InitState, Args))}
end).
```

```
F0 = newFoo(3).
```

```
\{V1, F1\} = F0(hi).
```

 $\{V2, F2\} = F1(6).$



...

Using a tuple for multiple methods (easily extend to record so methods have names)

```
newFoo(InitState) -> {
  fun (Args1) ->
     {f11(InitState, Args1),
     newFoo(f12(InitState, Args1))}
end,
fun (Args2) ->
     {f21(InitState, Args2),
     newFoo(f22(InitState, Args2))}
end,
```



Using Processes its clearer and easier

```
newServer(InitState) ->
Server = spawn(server),
{fun(Args1) ->
    rpc(Server, {f1, Args1}),
    fun(Args2) ->
    rpc(Server, {f2, Args2}),
    ...
}
end.
```

• rpc here is just the same rpc function we've seen before

f1 and f2 are atoms in this context



The server

```
server(State) ->
receive
{From, {f1, Args1}} ->
From ! f11(State, Args1),
server(f12(State, Args1));
{From, {f2, Args2}} ->
From ! f21(State, Args2),
server(f22(State, Args2))
```

end.

 f11, f12, f21, and f22 are the same functions as in the purely sequential simulation



The client

Server1 = newServer(SomeState),

- $\{$ S1F1, S1F2 $\}$ = Server1,
- V1 = S1F1(Args1),
- V2 = S1F2(Args2), % note V2 depends on both Args1 and Args2 as well as SomeState
- {S2F1, S2F2} = newServer(SomeOtherState),
- V3 = S2F1(OtherArgs1),
- V4 = S2F2(OtherArgs2), % note V4 depends
 on OtherArgs1 and 2 and SomeOtherState
 but not on SomeState, Args1 and Args2



The behavior is not quite the same

- If the sequential simulation is used from multiple processes each gets its own "fork" of the object history
 - Because the way it is used results in a new object every time a method is called
- The process simulation involves only a single object that, oh by the way, has synchronized access
 - A process, inherently, does only one thing at a time



A complete, simple example

```
-module(tester).
-export([start/1]).
start(Id) ->
  Server = spawn(fun () -> serverloop(ld, []) end),
  fun(Item) -> rpc( Server, Item ) end
serverloop(Id, History) ->
  receive
    { From, Something } ->
       io:format( "Server ~p: Got ~p from ~p with history ~p~n", [Id,
  Something, From, History]),
       From ! { self(), ok },
       serverloop(Id, [Something|History])
  end.
```