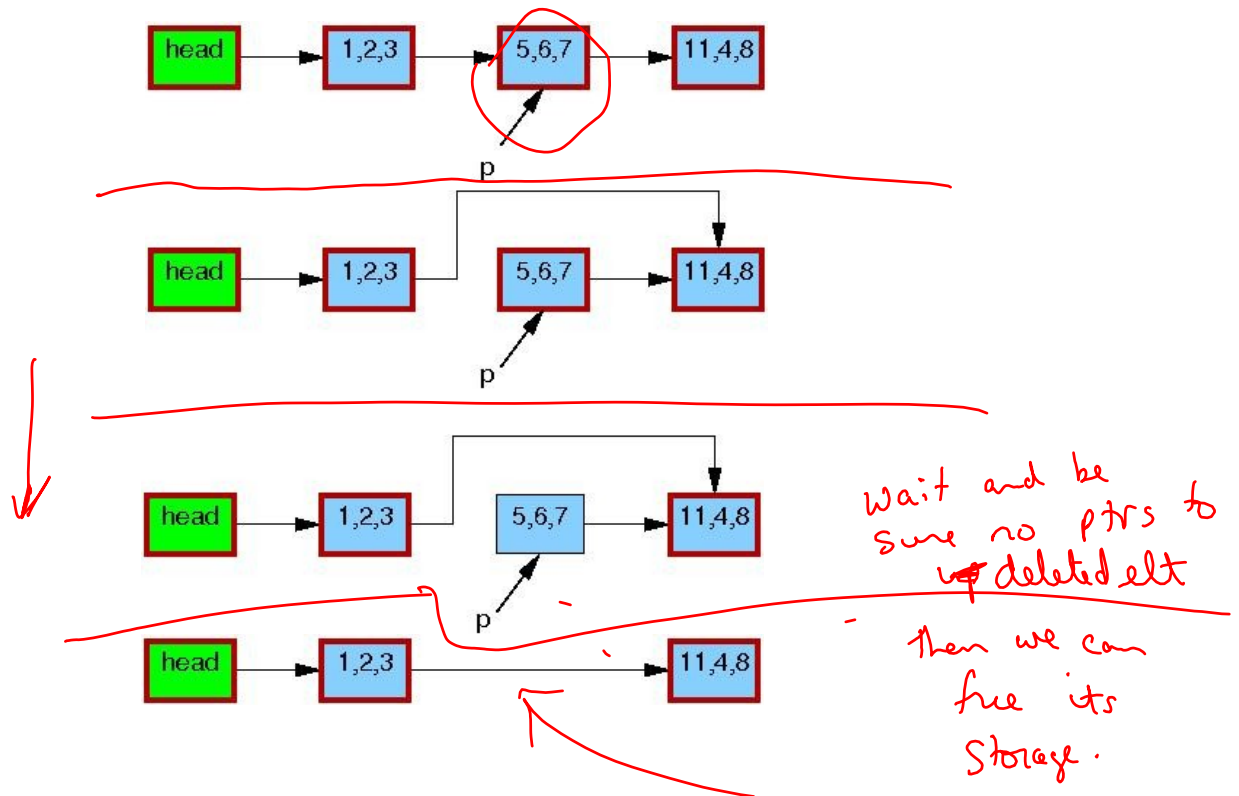


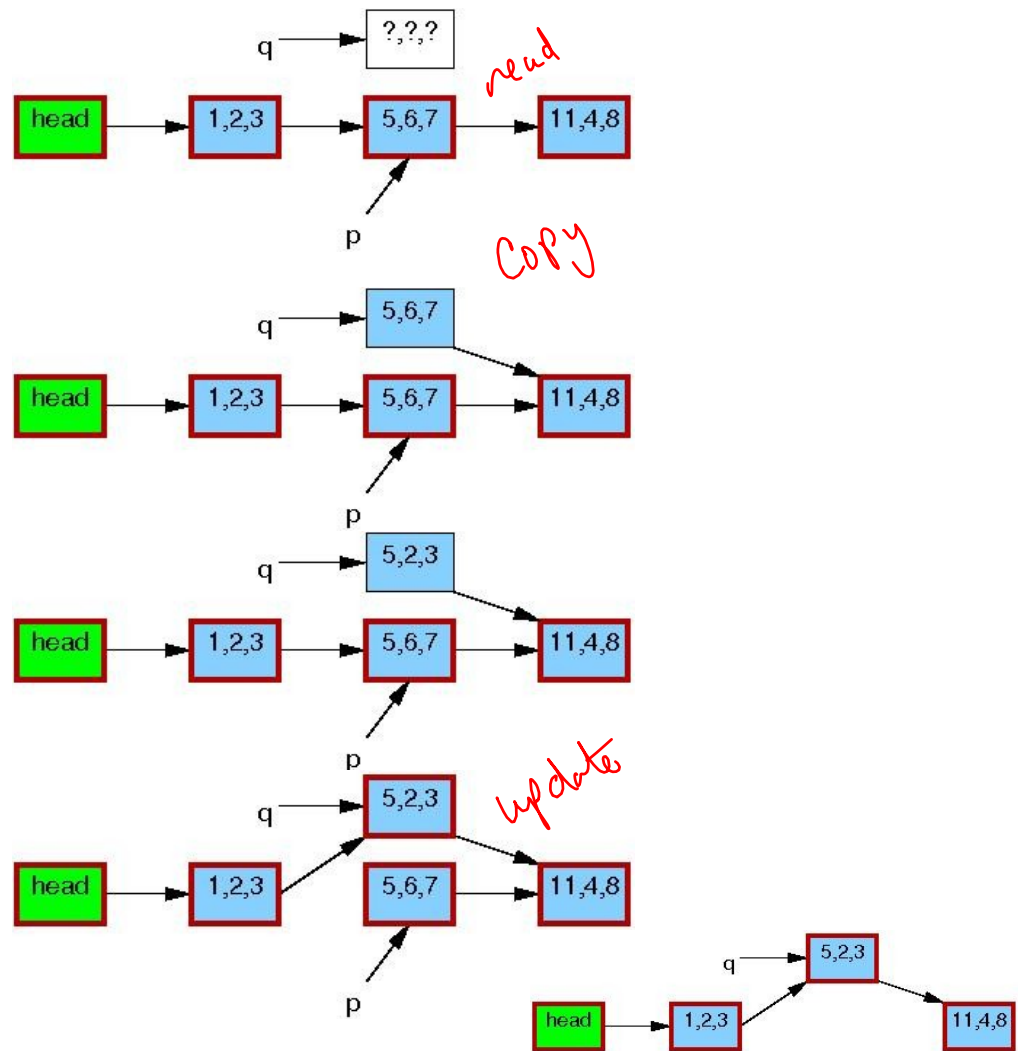
Read-Copy-Update (RCU)

CptS 483/580
April 21, 2010



Deleting a List Element

List element
update



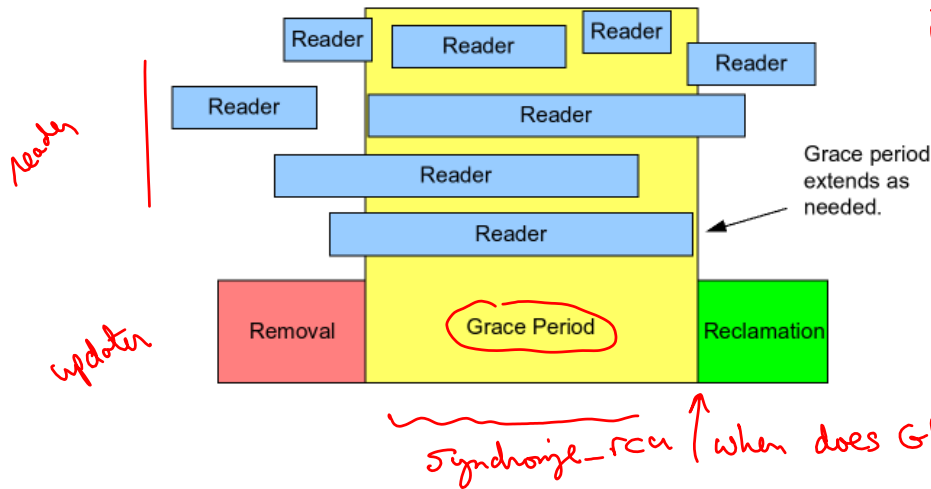
Read-copy-update: high-performance synchronization for read-mostly workloads.

Paul McKenney

Linux Weekly News (3 articles from late 2007-2008)

See <http://www.rdrop.com/users/paulmck/RCU/>

Figures here are from those three articles.



Note: update visibility is an issue that the Linux RCU interface also addresses with higher-level pointer read/update instructions that incorporate memory barriers

Primitives for RCU

```
rcu_read_lock(); /* read -side */  
rcu_read_unlock();  
synchronize_rcu(); /* write side */
```

before reading

done reading

returns when GP is known to be over.

synchronize_rcu() - implements the grace period after which removed elements may be freed.

Another approach is with a call-back procedure call_rcu(f,p) which calls f(p) when system determines that the grace period has expired.

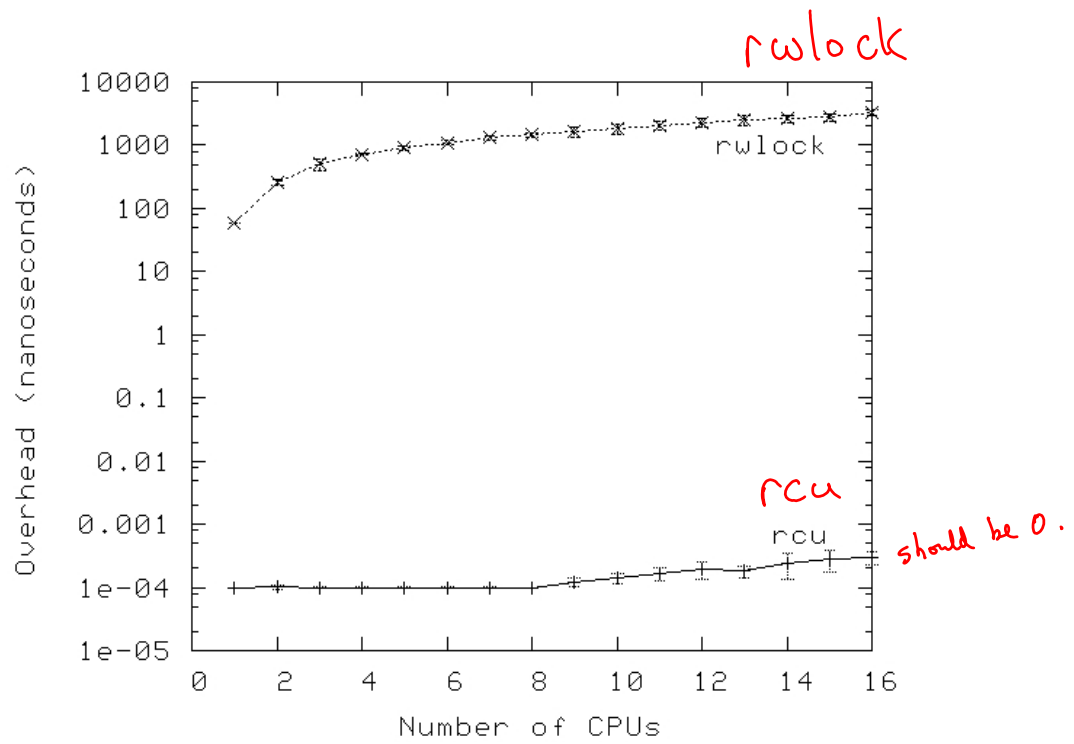
Implementing synchronize_rcu() – various mechanisms: if pre-emption is disallowed in lock...unlock sections sufficient to detect that all contexts have run once.

{ With garbage-collection: these techniques are (nearly) free – have to worry about update visibility.

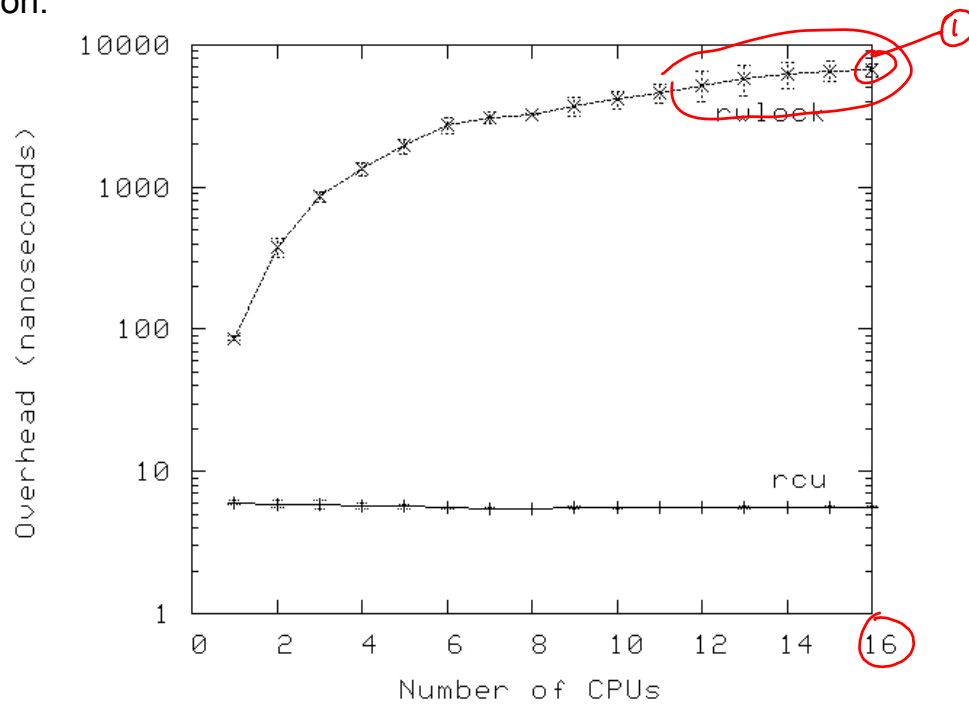
Note: locks (or TM, or ...) are still needed to protect against concurrent updates!

Micro-benchmarks

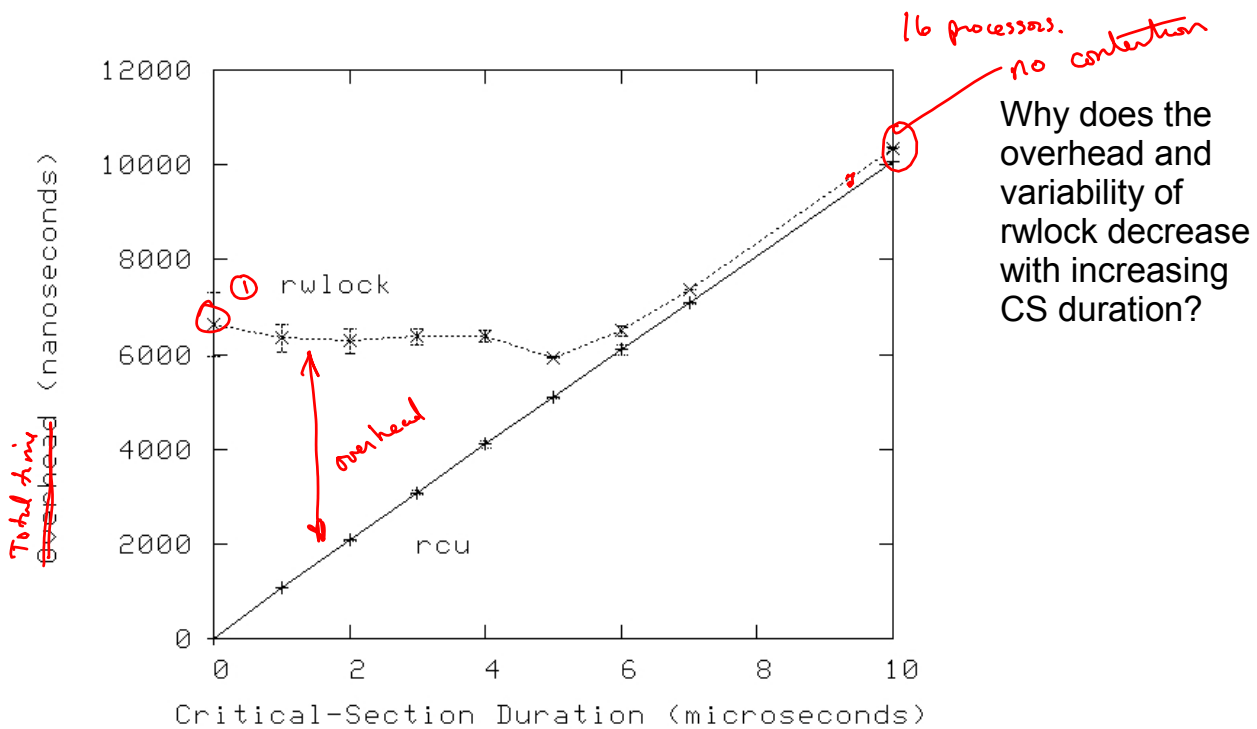
Overhead – non-preemptive kernel: 0



Overhead – preemptive kernel: non-zero cost because `read_lock_rcu()` has to turn off pre-emption.



Overhead: Pre-emptive kernel – 16
CPUs, increasing CS size



Macro-effect

RW-lock vs RCU

