

Lecture 22

Note Title

4/10/2008

Recall synch comm primitives

channels

send Evt

recv Evt

sync (single event)

select (list of events)

} potential for communication

} Project
Concurrent ML
CML

CML also has event combinators — functions that produce new events from existing events

wrap (e, f) when synchronized upon does the communication denoted by e and calls f on the result.

choose [e₁, e₂, ... e_n] is an event that synchronizes one of e₁, ... e_n

select [e₁, ... e_n] $\equiv$ sync(choose[e₁, ... e_n])

other combinators: $\swarrow$ ordinary value
(always Evt a) $\rightarrow$ evt that yields a when synchronized
never Evt $\rightarrow$ never successfully synchronizes
(timeout Evt t)

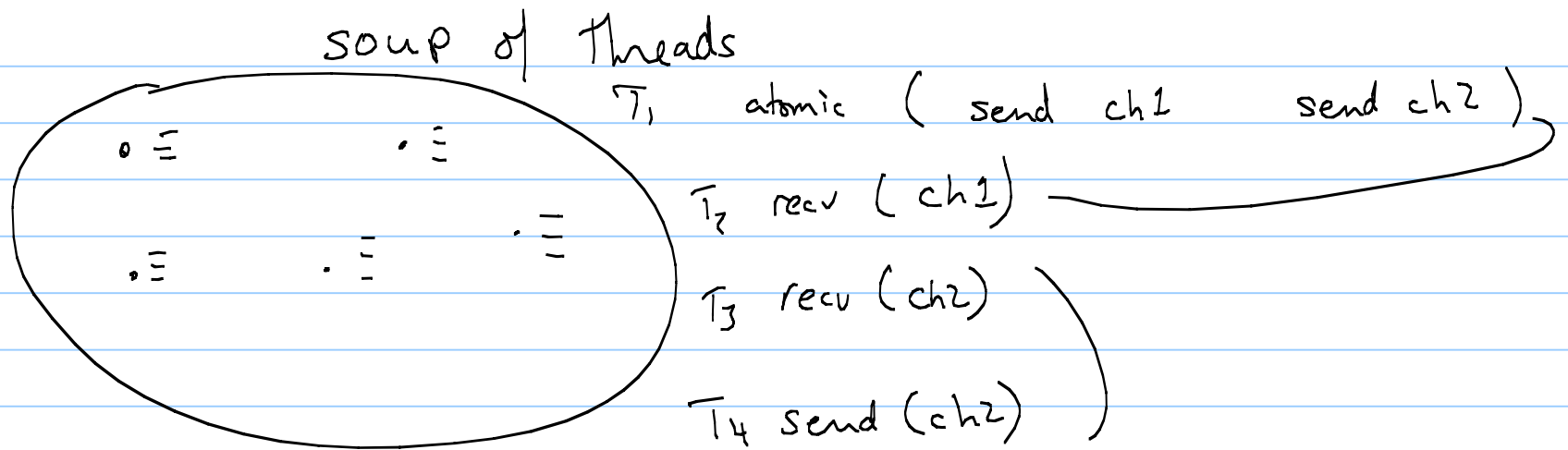
Fairly sophisticated comm. patterns can be built in event values — but not all interesting ones.

Consider

guarded receive — look at the value to be received and only take it if it is acceptable.

grecv Evt g $\swarrow$ boolean-valued function ch : only receives value x from ch if $(g\ x)$ is true.

Transactional Events - Fluet & Donnelly ICFP 2006 ~10 p.
" " Intl Journal of Functional Programming
2007 - 2008 - 79 pp



Power of the primitives: 1) emulate CML (not too surprising)
2) emulate STM

Implementation: currently based on STM

End of Synch Communication

locks, monitors, cv's, shared memory.

Asynchronous concurrent programming — doing concurrent i/o from a single thread — POSIX Asynch I/O interface, Win32 interface

— contrast to read, write syscalls.

— idea: prog. can launch multiple I/O ops; later find out the result by receiving an interrupt (signal) or by polling

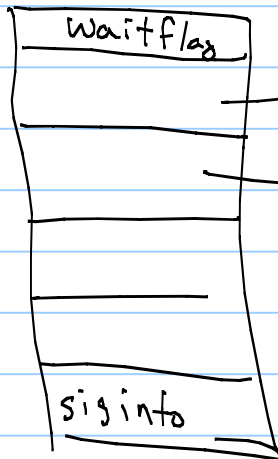
— poll/select system call model —

ask of a set of fds "are you ready to do i/o"
then when one is, you issue the read or write

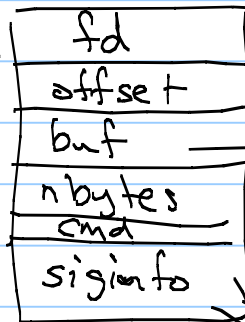
Posix.4:

Source: Programming for the Real World, Bill O. Gallmeister, O'Reilly.

listio



aio_cb



signal to deliver when this one is done

signal to deliver when all done

`aio_error (... acb ...)`
`aio_return (... acb ...)`

```
res = read ( ... )  
if (res < 0) {  
    e = errno
```

Signal handling: — C function set up as handler for a specific signal number.
passed the address of aio_cb assoc. w/
completed operation

```
if aio_error(acb) == EINPROGRESS { return }
```

```
if aio_return(acb) != acb->aio_nbytes {  
    maybe a problem
```

```
}
```

synchronization — protecting against concurrent access by a signal handler.

```
sigprocmask (SIG_BLOCK, & completion Signals, & prevMask)  
... do critical section ...
```

```
sigprocmask (SIG_SETMASK, & prevMask, NULL)
```

Why use AIO?

Use multiple resources concurrently

- other ways to accomplish this
 - threads (but sometimes implemented w/ AIO)
 - poll/select (also used to implement threads)

aio: start some i/o

→ find out when done
start some i/o

signals to identify being done
(painful)

poll/select:

wait for ready

do some i/o

wait for ready } compute for
a while then
wait for ready.

to program

nicer but not quite as
responsive