

CPTS580 In Class

Note Title

1/22/2008

fun {GateMaker F}

uses instantiation

fun {# Xs Ys}

uses genericity

fun {GateLoop Xs Ys}

case Xs # Ys of (X1 Xr) # (Y1 Yr) then

{F X Y} | {GateLoop Xr Yr} and

end

in thread {GateLoop Xs Ys} and

end

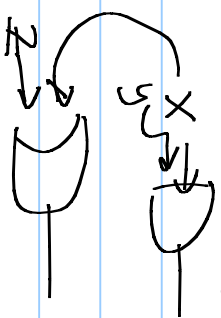
end

AndG = { GateMaker fun {# X Y} X * Y and }

OrG = { GateMaker " " X + Y - X * Y and }

AndG1 = { AndG Xs Ys }

OrG1 = { OrG Xs Zs }



Joining threads - synchronizing on termination of threads {and extracting a result value}.

local $x_1 \dots x_n$ in (C)

thread $\langle S \rangle_i$, $x_i = \underline{\text{unit}}$ end

thread $\langle S \rangle_2$ $x_2 = x_1$ and
:

thread $\langle S \rangle_n$ $x_n = x_{n-1}$ end

{ wait x_n }
end

local $x_1, \dots, x_n$ in

thread $\langle S \rangle_1$ $x_1 = \text{unit}$ end

:

thread $\langle S \rangle_n$ $x_n = \text{unit}$ end

{ wait x_1 } ... { wait x_n }

end

Is waiting necessary in the declarative model?

Coroutines

- a lot like a thread - it's a separate execution

context (stack, PC...) but switching betw.

co-routines is controlled by coroutines themselves
not by a scheduler.

$CID = \{ \text{Spawn } P \}$ - $\int$ 0-arg. procedure
if it does not
can at this point-
suspended.

$\{ \text{Resume } CID \}$ - $\{ P \}$

$\nearrow$ caller suspends order of execution

Python generators

Lazy evaluation (execution)

fun lazy $\{F A_1 \dots A_n\} \dots$ end normal function
Body

- $\{F A_1 \dots A_n\}$ creates a stopped execution
(a bit like a coroutine) that gets
activated when value of $\{F A_1 \dots A_n\}$ is needed.

to implement: augment kernel language with

By-need triggers.

$\{ByNeed P Y\}$ like thread $\{P Y\}$ end except

The thread is always executed, and the
trigger may not be. code variables that are triggered on.

implement $\{ByNeed <x> <y>\}$

if $<y>$ is determined (bound to a value)

Create a thread to evaluate $\{<x> <y>\}$

e.i. we add a trigger (x, y) (where x and
 y are the store variables associated to identifiers
 $<x>$ $<y>$ (in the current environment]) to the

trigger-shore.

← new component of an OS
Virtual machine.

Trigger is activated when a need for y is detected:

- a thread suspends waiting for y to be bound
 - a thread binds y .
- (remove trigger from trigger share)
- evaluation of $\{x < y\}$ takes place in a new thread.

| { }

The amazing thing is that B3Mod generates declarative needs.

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