

CptS 483/580 Feb. 26, 2013

Update Ring Assignment

No class this Thursday.

Mid term grades - based on exam and the early homeworks,

Implement shared state with a server process.

server(State) →

receive

{Pid, {get}} → Pid ! State, server(State);

{Pid, {set, NewState}} → Pid ! State, server(NewState);

{Pid, {update, UpdateFun}} → Pid ! State, ←
server(UpdateFun(State))

end.

init() → server([]).

So we're a user of this state server

register (myInt, spawn (fun server:init/0))).

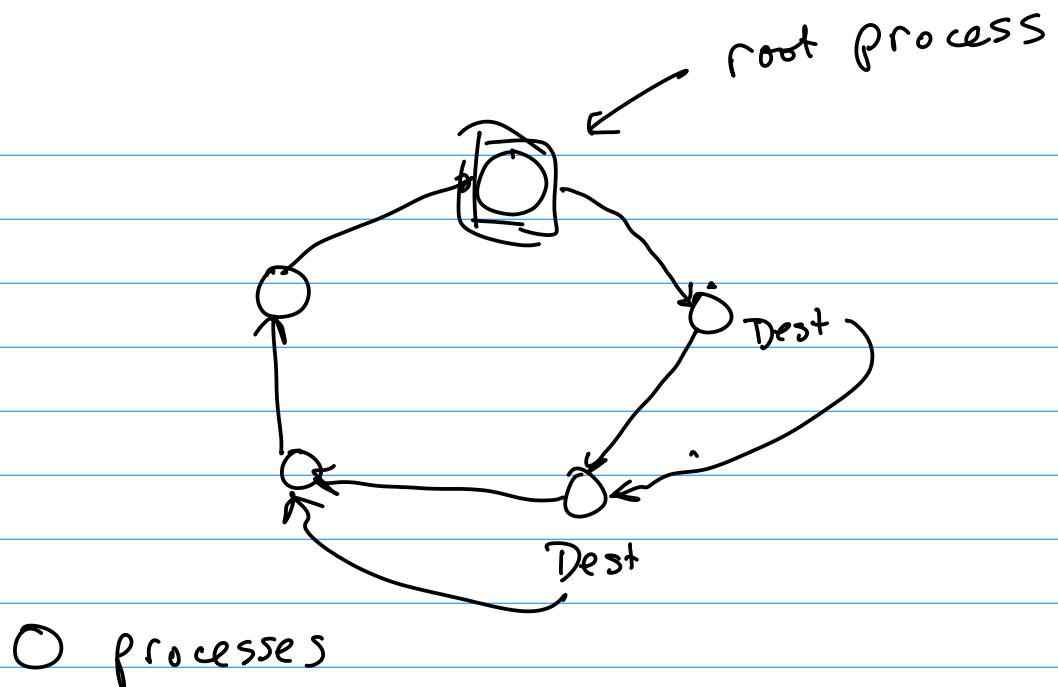
[myInt ! { self(), { set, 0 } },
receive
 Ans → Ans
end]

rpc (Server, Msg) →
 Server ! { self(), Msg },
receive
 Reply → Reply
end.

Ans = rpc (myInt, { set, 0 }).

get (Server) → rpc (Server, { get }).

set (Server, V) → rpc (Server, { set, V }).



○ processes

loop (Dest) →

receive
quit → Dest! quit;

M → Dest! M , loop (Dest)

end.

I - msg - at - a - time

Send All msgs

then

recv All msgs

Send some

→ recv what's available

Send some more

Suppose that's our implementation of the forwarder.

I have to build a lot of them and connect them together by writing in Dest differently in each one.

- { • each process somehow creates and initializes its successor?
- * • "master" process creates all of the forwarders and tells them to whom they should send. }

X

- {*} • Pass a parameter as part of spawning each process
• Send the dest as a message to each spawned process

}

→ `initter () →
receive
Dest → loop (Dest)
end.

param-batch-init(N) → param-batch-init-loop(N, self()).

param-batch-init-loop (N, First) →
lists:foldl (fun SpawnLoop/2, First, seg(1 N)).

SpawnLoop ($\frac{1}{A}$) → spawn (ring, loop, [A]).

What is the behavior of the root process?

Send some / receive some

mixed (M, N) \rightarrow

$Dest \leftarrow \text{param_batch_init}(N),$
 $\text{drive}(1, 1, M, Dest),$

$Dest ! \text{quit}.$

$\text{drive}(ToSend, ToRecv, Limit, Dest) \rightarrow$

$\text{Timeout} = \begin{cases} \text{if } ToSend < Limit \rightarrow 0; \\ \text{true} \rightarrow \text{infinity} \\ \text{end}, \end{cases}$

receive

$ToRecv$ when $ToRecv < Limit \rightarrow$

$\text{drive}(ToSend, ToRecv + 1, Limit, Dest);$

$ToRecv \rightarrow \{\}$

after Timeout \rightarrow

$Dest ! ToSend, \text{drive}(ToSend + 1, ToRecv, Limit, Dest),$
end.

A few hints:

use some output statements to make sure you are constructing the ring properly.

```
io=format ("Process ~p sends to ~p~n", [self(), Dest])
```

You don't want to do this in final version of your code -

In foot process, esp. if you are doing mixed - want to see interleaving of sending and receiving - use output statements in the receive and send branches of driver to print message numbers.

Again, not in the final code.